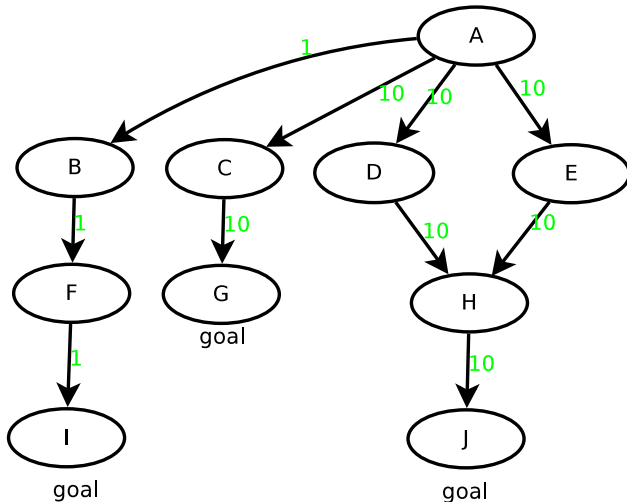


# CS444

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# Consider This Problem...



# Uniform Cost Search

**function** UNIFORM-COST-SEARCH(*problem*) **returns** a solution, or failure

*node* ← a node with STATE = *problem*.INITIAL-STATE, PATH-COST = 0

*frontier* ← a priority queue ordered by PATH-COST, with *node* as the only element

*explored* ← an empty set

**loop do**

**if** EMPTY?(*frontier*) **then return** failure

*node* ← POP(*frontier*) /\* chooses the lowest-cost node in *frontier* \*/

**if** *problem*.GOAL-TEST(*node*.STATE) **then return** SOLUTION(*node*)

  add *node*.STATE to *explored*

**for each** *action* **in** *problem*.ACTIONS(*node*.STATE) **do**

*child* ← CHILD-NODE(*problem*, *node*, *action*)

**if** *child*.STATE is not in *explored* or *frontier* **then**

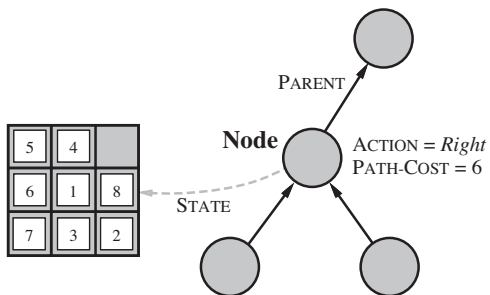
*frontier* ← INSERT(*child*, *frontier*)

**else if** *child*.STATE is in *frontier* with higher PATH-COST **then**

      replace that *frontier* node with *child*

**Figure 3.13** Uniform-cost search on a graph. The algorithm is identical to the general graph search algorithm in Figure ??, except for the use of a priority queue and the addition of an extra check in case a shorter path to a frontier state is discovered. The data structure for *frontier* needs to support efficient membership testing, so it should combine the capabilities of a priority queue and a hash table.

# Reminder... Search Nodes



# What We Really Want...

$$f(n) = g(n) + r(n)$$

Where,

$g(n)$  = cost to reach node  $n$ .

$r(n)$  = minimum cost to reach the goal, starting at  $n$ .

# What We Settle For...

$$f(n) = g(n) + h(n)$$

Where,

$g(n)$  = cost to reach node  $n$ .

$h(n)$  = *Estimate* of the minimum cost to reach the goal, starting at  $n$ .

Using  $f(n)$  to select nodes from the frontier gives us  $A^*$ .

- UCS
  - Complete
  - Optimal
- $A^*$ 
  - Complete
  - Optimal (If  $h(n)$  meets certain conditions)
  - Potentially much faster than UCS

- A\* is optimal for TREE-SEARCH if  $h(n)$  is admissible.
- A\* is optimal for GRAPH-SEARCH if  $h(n)$  is consistent.
- Admissible
  - $h(n)$  never overestimates the true cost to goal.
- Consistent
  - $h(n) \leq c(n, a, n') + h(n')$



# Examples...

# 8-Puzzle

7	2	4
5		6
8	3	1

Start State

	1	2
3	4	5
6	7	8

Goal State