



Graphs and Graph Models

Section 10.1

Graphs

Definition: A *graph* $G = (V, E)$ consists of a nonempty set V of *vertices* (or *nodes*) and a set E of *edges*. Each edge has either one or two vertices associated with it, called its *endpoints*. An edge is said to *connect* its endpoints.

Some Terminology

- In a *simple graph* each edge connects two different vertices and no two edges connect the same pair of vertices.
- *Multigraphs* may have multiple edges connecting the same two vertices. When m different edges connect the vertices u and v , we say that $\{u,v\}$ is an edge of *multiplicity* m .
- An edge that connects a vertex to itself is called a *loop*.
- A *pseudograph* may include loops, as well as multiple edges connecting the same pair of vertices.

Directed Graphs

Definition: An *directed graph* (or *digraph*) $G = (V, E)$ consists of a nonempty set V of *vertices* (or *nodes*) and a set E of *directed edges* (or *arcs*). Each edge is associated with an ordered pair of vertices. The directed edge associated with the ordered pair (u, v) is said to *start at* u and *end at* v .

Remark:

- Graphs where the end points of an edge are not ordered are said to be *undirected graphs*.

Graph Terminology: Summary

TABLE 1 Graph Terminology.			
<i>Type</i>	<i>Edges</i>	<i>Multiple Edges Allowed?</i>	<i>Loops Allowed?</i>
Simple graph	Undirected	No	No
Multigraph	Undirected	Yes	No
Pseudograph	Undirected	Yes	Yes
Simple directed graph	Directed	No	No
Directed multigraph	Directed	Yes	Yes
Mixed graph	Directed and undirected	Yes	Yes



Graph Terminology and Special Types of Graphs

Section 10.2

Basic Terminology

Definition 1. Two vertices u, v in an undirected graph G are called *adjacent* (or *neighbors*) in G if there is an edge e between u and v . Such an edge e is called *incident with* the vertices u and v and e is said to *connect* u and v .

Definition 2. The set of all neighbors of a vertex v of $G = (V, E)$, denoted by $N(v)$, is called the *neighborhood* of v . If A is a subset of V , we denote by $N(A)$ the set of all vertices in G that are adjacent to at least one vertex in A . So,

$$N(A) = \bigcup_{v \in A} N(v).$$

Definition 3. The *degree of a vertex in a undirected graph* is the number of edges incident with it, except that a loop at a vertex contributes two to the degree of that vertex. The degree of the vertex v is denoted by $\deg(v)$.

Degrees of Vertices

● **Theorem 1 (*Handshaking Theorem*):** If $G = (V, E)$ is an undirected graph with m edges, then

$$2m = \sum_{v \in V} \deg(v)$$

Proof:

Each edge contributes twice to the degree count of all vertices. Hence, both the left-hand and right-hand sides of this equation equal twice the number of edges. ▶

Think about the graph where vertices represent the people at a party and an edge connects two people who have shaken hands.

Directed Graphs

Recall the definition of a directed graph.

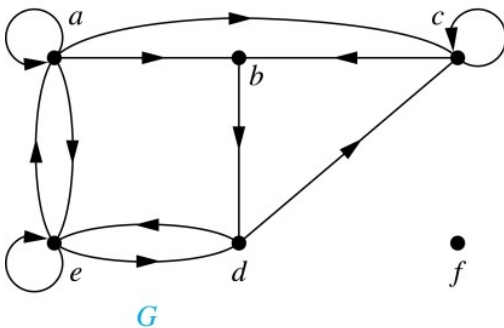
Definition: An *directed graph* $G = (V, E)$ consists of V , a nonempty set of *vertices* (or *nodes*), and E , a set of *directed edges* or *arcs*. Each edge is an ordered pair of vertices. The directed edge (u, v) is said to start at u and end at v .

Definition: Let (u, v) be an edge in G . Then u is the *initial vertex* of this edge and is *adjacent to* v and v is the *terminal* (or *end*) *vertex* of this edge and is *adjacent from* u . The initial and terminal vertices of a loop are the same.

Directed Graphs (*continued*)

Definition: The *in-degree* of a vertex v , denoted $\deg^-(v)$, is the number of edges which terminate at v . The *out-degree* of v , denoted $\deg^+(v)$, is the number of edges with v as their initial vertex.

Example: In the graph G we have



$$\deg^-(a) = 2, \deg^-(b) = 2, \deg^-(c) = 3, \\ \deg^-(d) = 2, \deg^-(e) = 3, \deg^-(f) = 0.$$

$$\deg^+(a) = 4, \deg^+(b) = 1, \deg^+(c) = 2, \\ \deg^+(d) = 2, \deg^+(e) = 3, \deg^+(f) = 0.$$

Directed Graphs (*continued*)

Theorem 3: Let $G = (V, E)$ be a graph with directed edges. Then:

$$|E| = \sum_{v \in V} \text{deg}^{-}(v) = \sum_{v \in V} \text{deg}^{+}(v).$$


Proof: The first sum counts the number of outgoing edges over all vertices and the second sum counts the number of incoming edges over all vertices. It follows that both sums equal the number of edges in the graph. ▶

Special Types of Simple Graphs: Complete Graphs

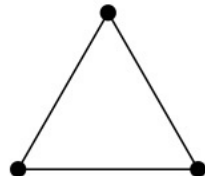
A *complete graph on n vertices*, denoted by K_n , is the simple graph that contains exactly one edge between each pair of distinct vertices.



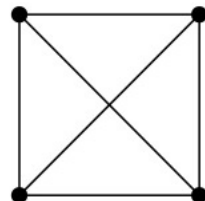
K_1



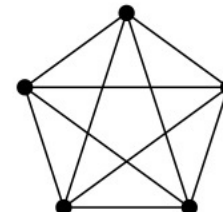
K_2



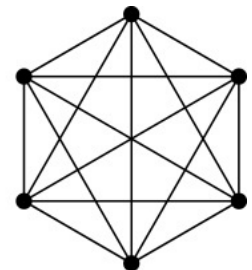
K_3



K_4



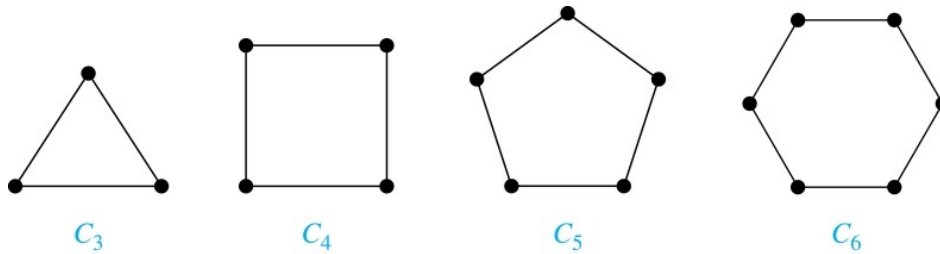
K_5



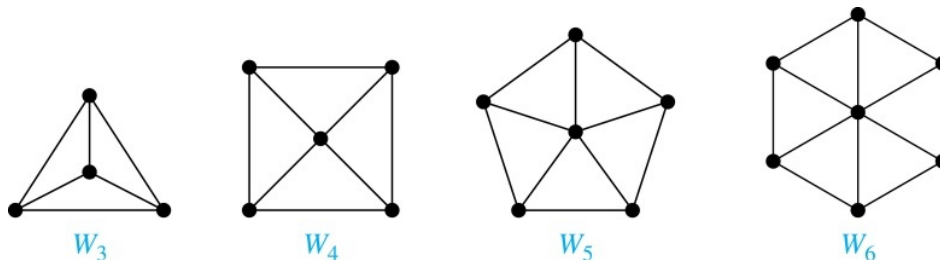
K_6

Special Types of Simple Graphs: Cycles and Wheels

A *cycle* C_n for $n \geq 3$ consists of n vertices v_1, v_2, \dots, v_n , and edges $\{v_1, v_2\}, \{v_2, v_3\}, \dots, \{v_{n-1}, v_n\}, \{v_n, v_1\}$.

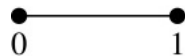


A *wheel* W_n is obtained by adding an additional vertex to a cycle C_n for $n \geq 3$ and connecting this new vertex to each of the n vertices in C_n by new edges.

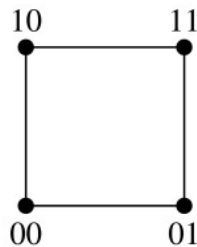


Special Types of Simple Graphs: n -Cubes

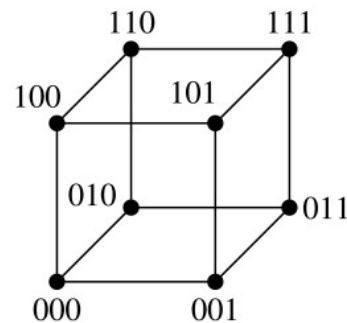
An n -dimensional hypercube, or n -cube, Q_n , is a graph with 2^n vertices representing all bit strings of length n , where there is an edge between two vertices that differ in exactly one bit position.



Q_1



Q_2



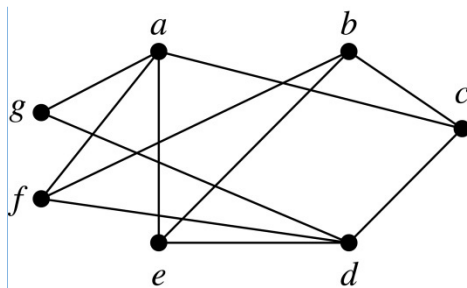
Q_3

Bipartite Graphs

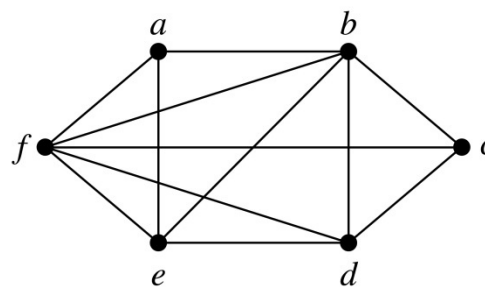
Definition: A simple graph G is bipartite if V can be partitioned into two disjoint subsets V_1 and V_2 such that every edge connects a vertex in V_1 and a vertex in V_2 . In other words, there are no edges which connect two vertices in V_1 or in V_2 .

It is not hard to show that an equivalent definition of a bipartite graph is a graph where it is possible to color the vertices red or blue so that no two adjacent vertices are the same color.

G is bipartite



G

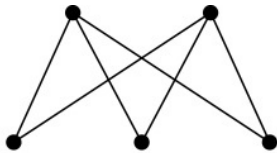


H

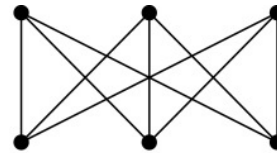
H is not bipartite since if we color a red, then the adjacent vertices f and b must both be blue.

Complete Bipartite Graphs

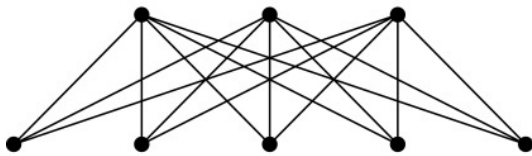
Definition: A complete bipartite graph $K_{m,n}$ is a graph that has its vertex set partitioned into two subsets V_1 of size m and V_2 of size n such that there is an edge from every vertex in V_1 to every vertex in V_2 .



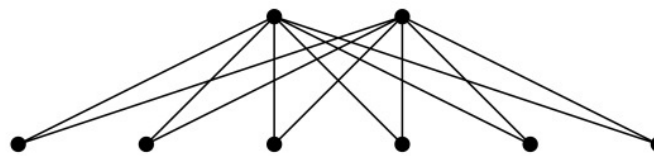
$K_{2,3}$



$K_{3,3}$



$K_{3,5}$

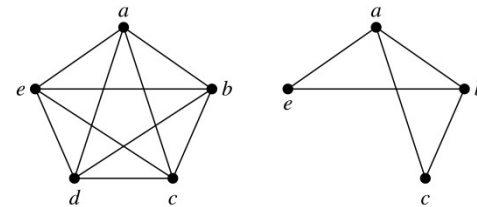


$K_{2,6}$

New Graphs from Old

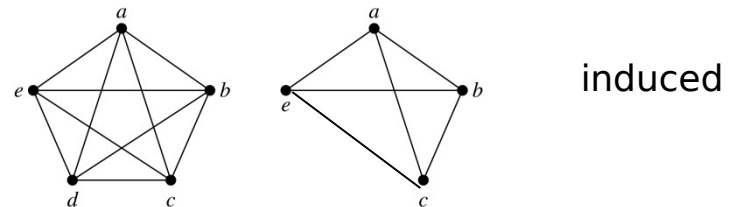
Definition: A subgraph of a graph $G = (V, E)$ is a graph (W, F) , where $W \subset V$ and $F \subset E$. A subgraph H of G is a proper subgraph of G if $H \neq G$.

Example: Here we show K_5 and one of its subgraphs.



Definition: Let $G = (V, E)$ be a simple graph. The *subgraph induced* by a subset W of the vertex set V is the graph (W, F) , where the edge set F contains an edge in E if and only if both endpoints are in W .

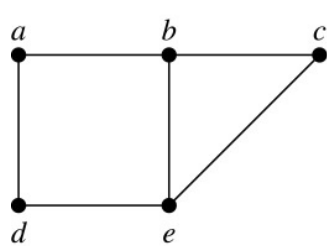
Example: Here we show K_5 and the subgraph by $W = \{a, b, c, e\}$.



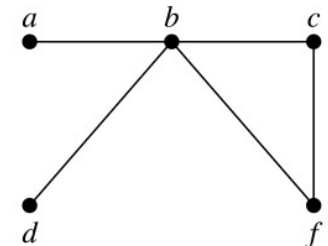
New Graphs from Old (*continued*)

Definition: The *union* of two simple graphs $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ is the simple graph with vertex set $V_1 \cup V_2$ and edge set $E_1 \cup E_2$. The union of G_1 and G_2 is denoted by $G_1 \cup G_2$.

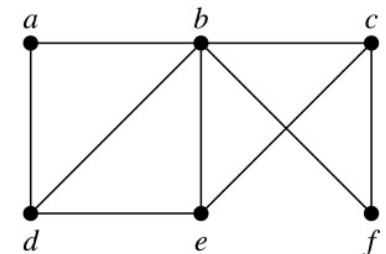
Example:



G_1



G_2



$G_1 \cup G_2$

(a)

(b)