# Artificial Intelligence

# Logical Agents and Propositional Logic

CS 444 – Spring 2020
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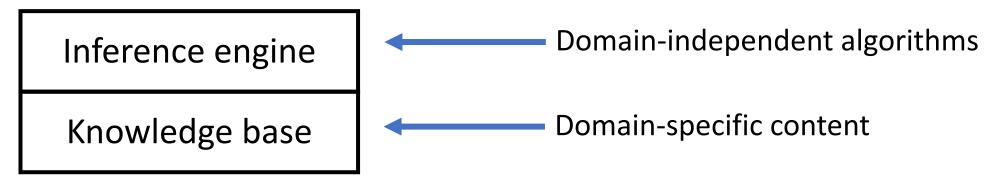


## Outline

- Knowledge-based Agents
- Wumpus World
- Logic Models and Entailment
- Propositional (Boolean Logic)
- Model checking: Inference by Enumeration



# Knowledge Bases



Knowledge base = set of sentences in a formal language Declarative approach to build an agent (or other system):

- Tell it what it needs to know
- Then it can Ask itself what to do, answers should follow from the KB

Agents can be viewed at the knowledge level.

i.e., what they know, regardless of how implemented

Or at the implementation level

i.e., data structures in KB and algorithms that manipulate them



# A Simple Knowledge-based Agent

function KB-Agent(percept) returns an action

Static: KB, a knowledge base

t, a counter, initially 0, that indicates time

Tell(KB, Make-Percept-Sentence(percept, t))

action ← Ask(KB, Make-Action-Query(t))

Tell(KB, Make-Percept-Sentence(action, t))

 $t \leftarrow t + 1$ 

return action

#### The agent must be able to:

- Represent states, actions, etc.
- Incorporate new percepts
- Update internal representations of the world
- Deduce hidden properties of the world
- Deduce appropriate actions



# Wumpus World – PEAS Description

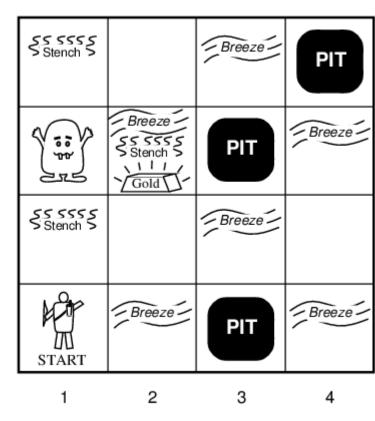
#### **Performance measure**

gold + 1000, death – 1000

-1 per step, -10 for using the arrow

#### **Environment**

- Squares adjacent to Wumpus are smelly
- Squares adjacent to pit are breezy
- glitter iff gold is in the same square
- Shooting kills Wumpus if you are facing it
- Shooting uses up the only arrow
- Grabbing picks up gold if in same square
- Releasing drops the gold in same square
- Squeal sound means Wumpus was killed



#### **Actuators**

Left turn, right turn, forward, grab, release, shoot

#### Sensors

Breeze, glitter, smell

3

2



# Wumpus World Characterizations

Observable? Partially observable – only local perception

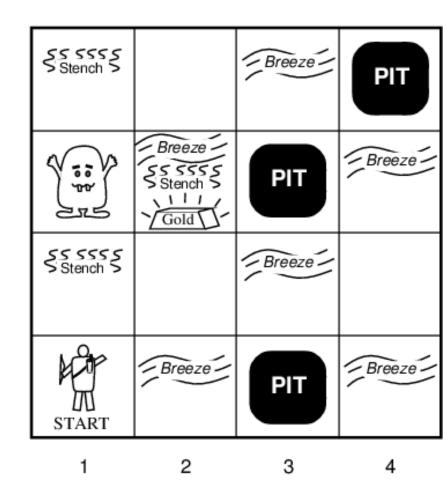
Deterministic? Yes – outcomes exactly specified

Episodic? No – sequential at the level of actions

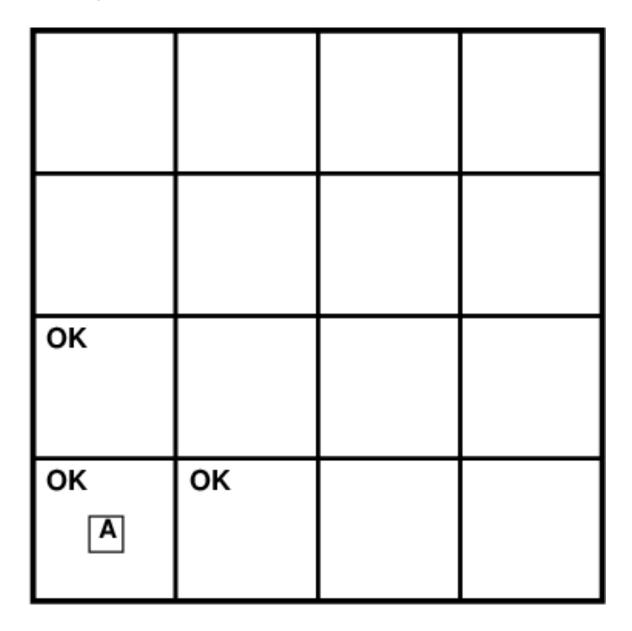
Static? Yes – Wumpus and pits do not move

Discrete? Yes

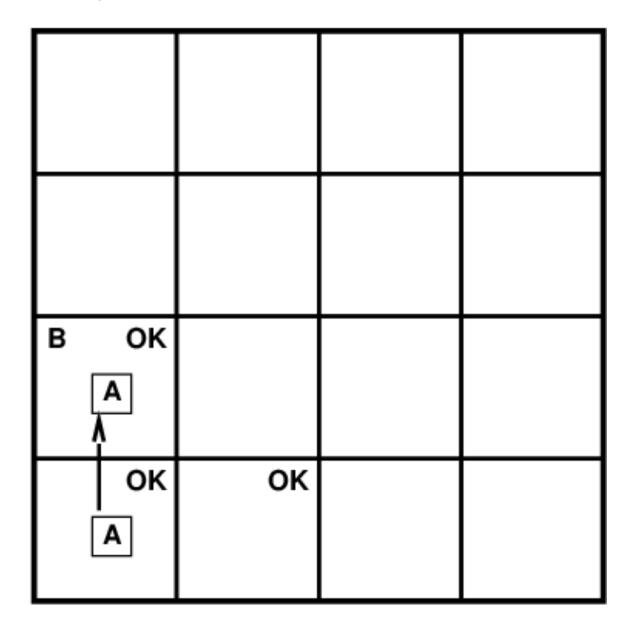
Single Agent? Yes (Wumpus is essentially a natural feature)



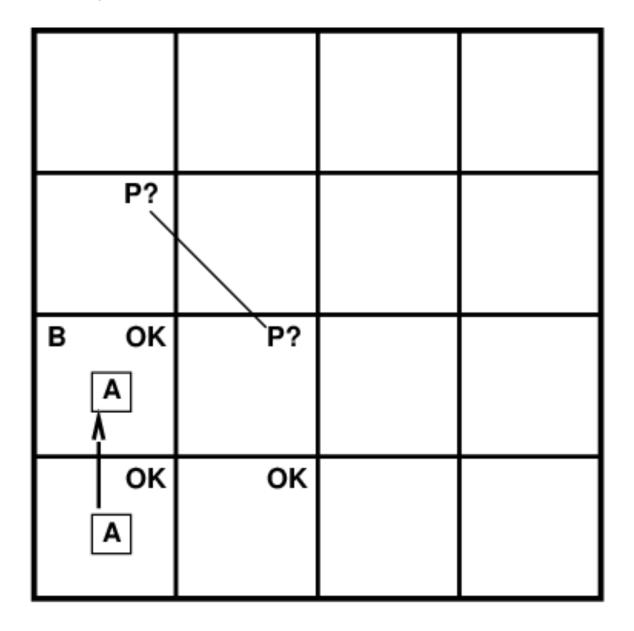




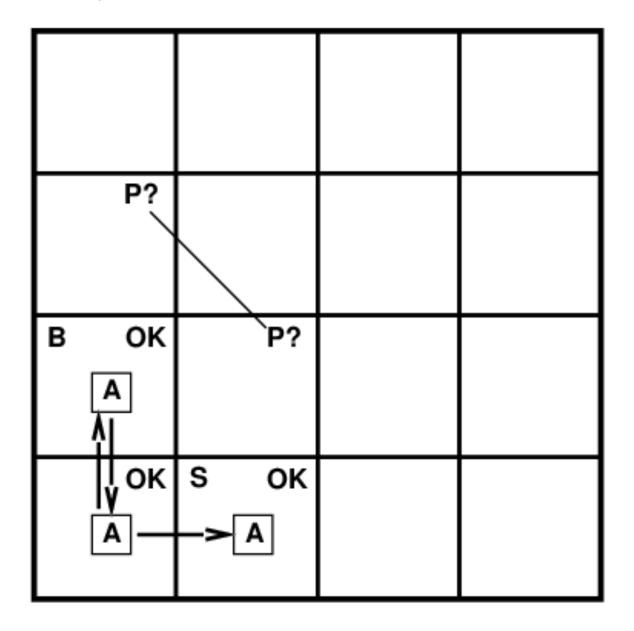




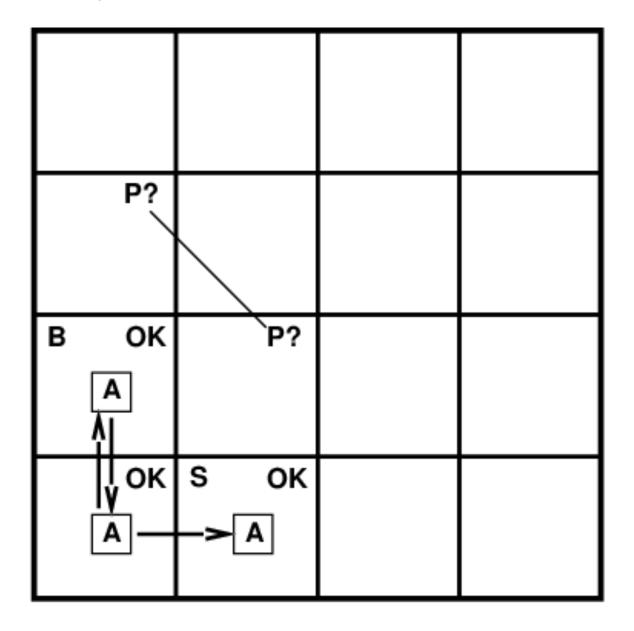




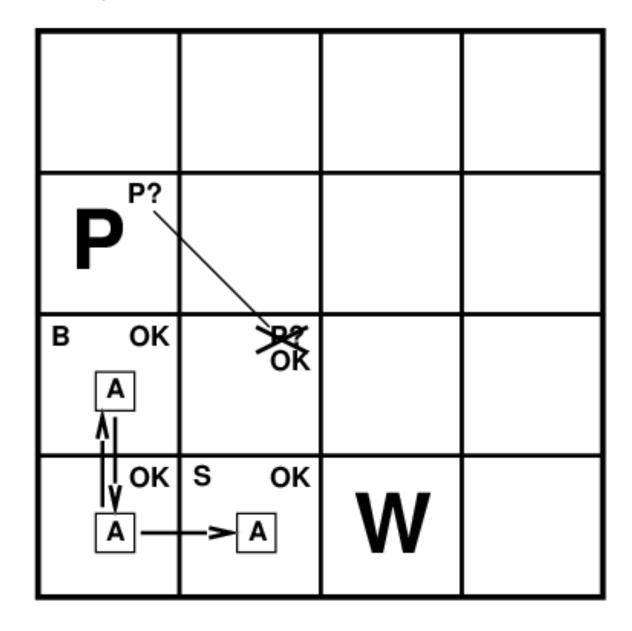




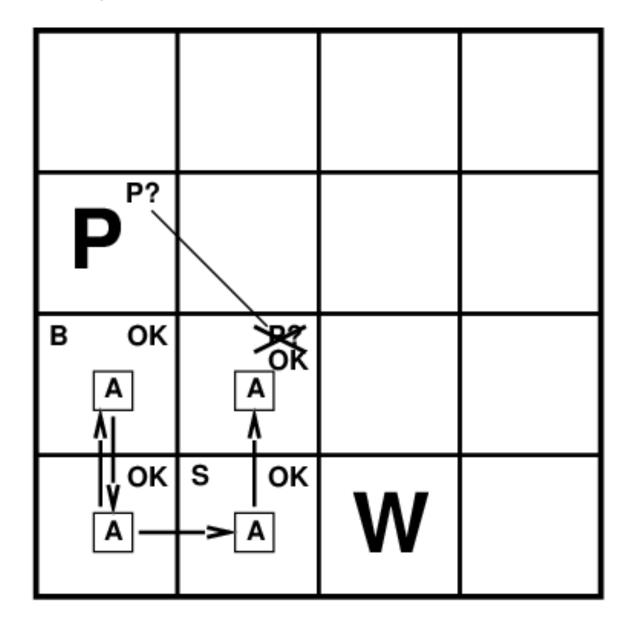




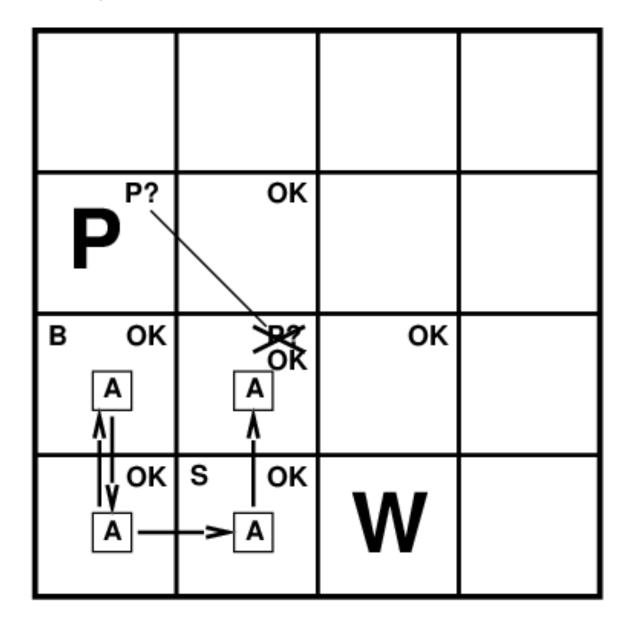




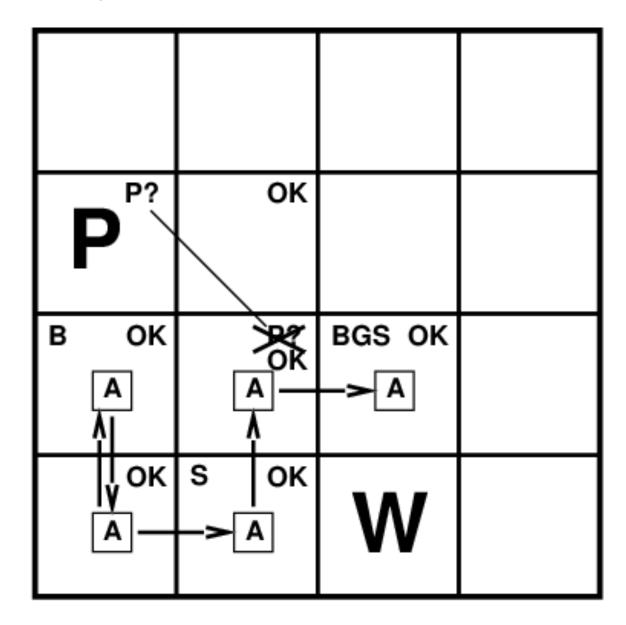






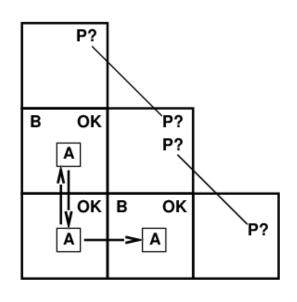








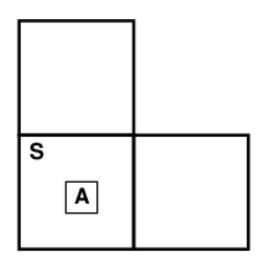
#### Sometimes No Safe Move Exists



Breeze in (1,2) and (2, 1)

⇒no safe actions

Assuming pits uniformly distributed, (2,2) has a pit with higher probability. How much?



Smell in (1, 1)

⇒cannot move

Can use strategy of coercion:

shoot straight ahead

Wumpus was there ⇒ dead ⇒ safe

Wumpus wasn't there ⇒ safe



#### Logic in General

- Logics are formal languages for representing information such that conclusions can be drawn.
- Syntax determines how sentences are expressed in a particular logic/language
- Semantics define the "meaning of sentences";
  - i.e., define truth of a sentence in a world
- e.g., the language of arithmetic:
  - $x + 2 \ge y$  is a sentence;  $x^2 + y > 1$  is not a sentence
  - $x + 2 \ge y$  is true iff the number x + 2 is no less than the number y
  - $x + 2 \ge y$  is true in a world where x = 7, y = 1
  - $x + 2 \ge y$  is false in a world where x = 0, y = 6



## Types of Logic

- Logics are characterized by what they commit to as primitives
- Ontological commitment: what exists facts? Objects? Time? Beliefs?
- Epistemological commitment: what states of knowledge?

Language	Ontological Commitment	Epistemological Commitment
Propositional logic	Facts	true/false/unknown
First-order logic	Facts, objects, relations	true/false/unknown
Temporal logic	Facts, objects, relations, time	true/false/unknown
Probability theory	Facts	Degree of belief
Fuzzy logic	Facts + degree of truth	Known internal value



## Reasoning with Logic

- First order of business: fundamental concepts of logic representation. And reasoning
  - independent of any logic's particular form/type
  - Entailment
- Second order of business: Introduction to propositional logic
   Wumpus KB via propositional logic
- Third order of business: Drawing conclusions
  - Inference and theorem proving



#### Models

- Can use the term model in place of possible world
- Logicians typically think in terms of models, which are formally-structed worlds with respect to which truth can be evaluated
- Model = mathematical abstraction that fixes the truth/falsehood of every relevant sentence
- Possible models are just all possible assignments of variables in the environment
- We say that a model m "satisfies" sentence  $\alpha$  if  $\alpha$  "is true in" m Or: "m is a model of  $\alpha$ "  $M(\alpha)$  is the set of all models of  $\alpha$



## Models and Entailment

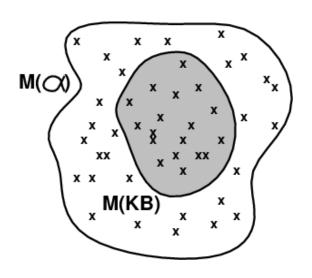
**Entailment** means that one thing follows from another:

$$KB \models \alpha$$

Knowledge base KB entails sentence  $\alpha$ 

iff  $\alpha$  is true in all worlds/models where KB is true

$$KB \models \alpha \text{ iff } M(KB) \subseteq M(\alpha)$$



e.g., KB contains "Giants won" and "Red won" entails "Giants or Reds won"

$$x + y = 4$$
 entails  $4 = x + y$ 

Entailment is a relationship between sentences (i.e., syntax) that is based on semantics.

Note: brains process syntax (or some sort).



#### Quick Exercise

Given two sentences  $\alpha$  and  $\beta$ , what does this mean:

$$\alpha \vDash \beta$$

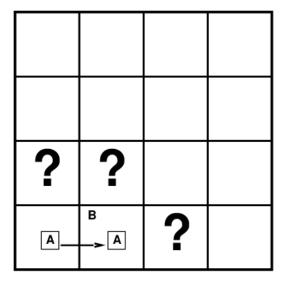
- $\alpha$  entails  $\beta$
- $M(\alpha) \subseteq M(\beta)$
- $\beta$  is satisfied in all models of  $\alpha$ .
- $\beta$  may be satisfied in other models (as well).
- $\alpha$  is stronger assertion than  $\beta$ .

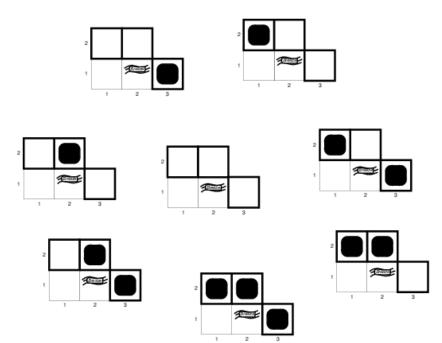
## Hands on: Entailment in the Wumpus World

Situation after detecting nothing in [1, 1], moving right, breeze in [2,1]

Consider possible models for ?s assuming only pits

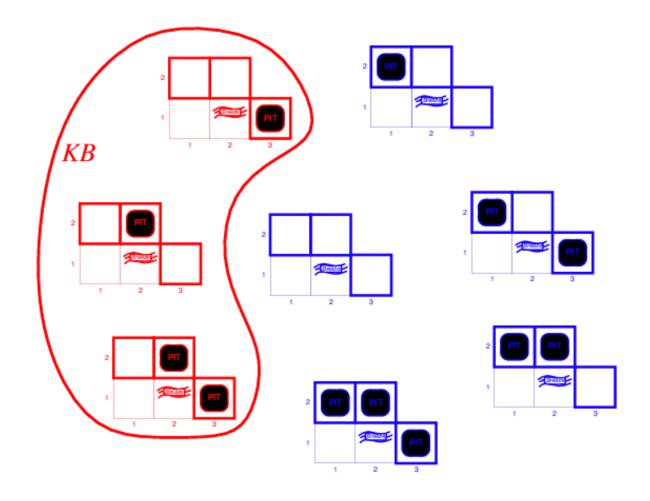
3 Boolean choices ⇒ 8 possible models







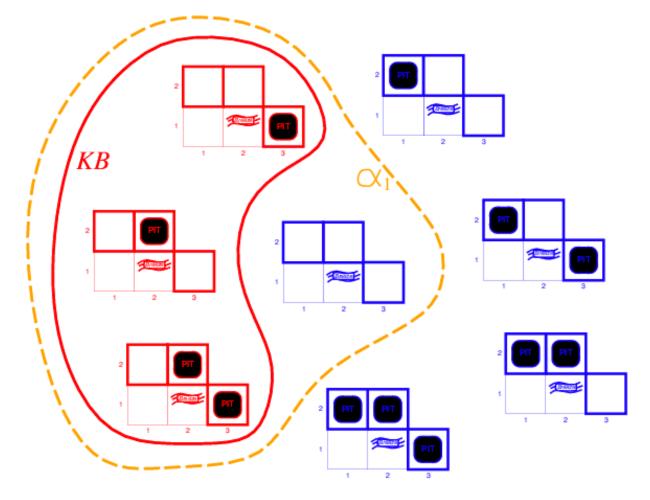
## Wumpus Models



**KB** = Wumpus-world rules + observations



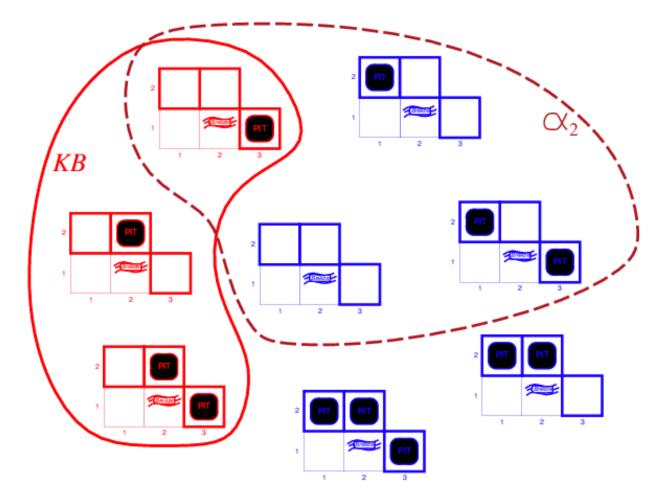
#### Wumpus Models



KB = Wumpus-world rules + observations,  $\alpha_1$  = "[1,2] is safe" KB  $\models \alpha_1$ , proved by model checking



#### Wumpus Models



KB = Wumpus-world rules + observations,  $\alpha_2$  = "[2,2] is safe", KB  $\not\models \alpha_2$ 



# From Entailment to Logical Inference

Entailment can be used to derive logical conclusions

i.e., carry out logical inference

A straightforward algorithm to carry out inference:

#### **Model Checking**

Model checking enumerates all possible models to check that  $\alpha$  is true in all models where KB is true. i.e., M(KB)  $\subseteq$  M( $\alpha$ )

To understand entailment and inference: haystack and needle analogy

Consequences of KB are a haystack,  $\alpha$  is a needle.

Entailment = needle in haystack

Inference = finding it

We need inference procedures to derive  $\alpha$  from a given KB.



#### Inference

 $KB \vdash_i \alpha = \text{sentence } \alpha \text{ can be derived from } KB \text{ by procedure } i$ 

Soundness: inference procedure i is sound if whenever  $KB \vdash_i \alpha$ , it is also true that  $KB \models \alpha$  (does not make stuff up)

Completeness: inference procedure i is complete if whenever  $KB \models \alpha$ , it is also true that  $KB \vdash_i \alpha$  (finds the needle in haystack)

Preview: we will define a logic (first-order logic) which is expressive enough to say almost anything of interest, and for which there exists a sound and complete inference procedure.

That is, the procedure will answer any question whose answer follows from what is known by the KB.

First step, propositional logic.



# Propositional Logic: Syntax

Propositional logic is the simplest logic – illustrates basic ideas

Atomic sentences consist of a single propositional symbol e.g., Propositional symbols  $P_1$ ,  $P_2$ , etc. are atomic sentences

Each such symbol stands for a proposition that can be true or false. e.g., W1,3 stands for proposition that Wumpus is in [1,3]

Two propositions with fixed meaning: True and False

Complex sentences build over atomic ones via connectives: negation, conjunction, disjunction, implication, biconditional



# Propositional Logic: Syntax

```
If S is a sentence, \neg S is a sentence (negation)
      A (positive) literal is an atomic sentence; a (negative) literal is a negated atomic sentence
If S_1 and S_2 are sentences S_1 \wedge S_2 is a sentence (conjunction)
  S<sub>1</sub> and S<sub>2</sub> are called conjuncts
If S_1 and S_2 are sentences S_1 \vee S_2 is a sentence (disjunction)
  S<sub>1</sub> and S<sub>2</sub> are called disjuncts
If S_1 and S_2 are sentences, S_1 \Longrightarrow S_2 is a sentence (implication/conditional)
    S<sub>1</sub> is called premise/antecedent; S<sub>2</sub> is called conclusion or consequent
    Implication also known as rule or if-then statement
If S_1 and S_2 are sentences, S_1 \Leftrightarrow S_2 is a sentence (biconditional)
```



## Propositional Logic: Semantics (Backus-Naur Form (BNF)

BNF is an ambiguous formal grammar for propositional logic

Sentence → AtomticSentence | ComplexSentence

AtomicSentence → True | False | P | Q | ...

ComplexSentence → (Sentence) | [Sentence] | ¬Sentence | Sentence ∧ Sentence ...

We add operator precedence to disambiguate

Operator precedence (from highest to lowest):  $\neg$ ,  $\land$ ,  $\lor$ ,  $\Longrightarrow$ ,  $\Leftrightarrow$ 



## Propositional Logic: Semantics

Each model specifies true/false for each propositional symbol

e.g., P<sub>1,2</sub> true

P<sub>2,2</sub> true P<sub>3,1</sub> false

This specific model:  $m_1 = \{P_{1,2} = \text{true}, P_{2,2} = \text{true}, P_{3,1} = \text{false}\}$ 

 $2^3$  = 8 possible models, feasible to enumerate

Rules for evaluating truth with respect to a model m:

¬S	is true iff	S	S is false		
$S_1 \wedge S_2$	is true iff	S <sub>1</sub>	is true and	S <sub>2</sub>	is true
$S_1 \vee S_2$	is true iff	S <sub>1</sub>	is true or	S <sub>2</sub>	is true
$S_1 \Longrightarrow S_2$	is true iff	S <sub>1</sub>	is false or	S <sub>2</sub>	is true
i.e.	is false iff	S <sub>1</sub>	is true and	S <sub>2</sub>	is false
$S_1 \Leftrightarrow S_2$	is true iff	$S_1 \Longrightarrow S_2$	is true and	$S_2 \Longrightarrow S_1$	is true
$_{2} \wedge (P_{2,2} \vee p_{3,1})$	= true ∧ (	false V tr	ue) = tr	ue∧true	= true



#### Truth Table for Connectives

P	Q	¬P	PΛQ	PVQ	$P \Longrightarrow Q$	$P \Leftrightarrow Q$
false	false	true	false	false	true	true
false	true	true	false	true	true	false
true	false	false	false	true	false	false
true	true	false	true	true	true	true



## Wumpus World Sentences in Propositional Logic

Let P<sub>i,i</sub> be true if there is a pit in [i, j]

Let B<sub>i,i</sub> be true if agent is in [i, j] and perceives a breeze

Let W<sub>i,j</sub> be true if there is a Wumpus in [i, j]

Let S<sub>i,i</sub> be true if agent is in [i, j] and perceives a stench

... you can define other atomic sentences

Percept sentences part of KB:

No pit, no breeze in [1, 1], but breeze perceived when in [2, 1]

$$R_1 : \neg P_{1,1}; \quad R_4 : \neg B_{1,1} \qquad \qquad R_5 : B_{2,1}$$

Rules in KB: "Put cause breezes in adjacent squares" eqv. to "Square is breeze iff adjacent pit".

$$R_2: B_{1,1} \Leftrightarrow (P_{1,2} \vee P_{2,1})$$

$$R_3: B_{2,1} \Leftrightarrow (P_{1,1} \vee P_{2,2} \vee P_{3,1})$$



## Truth Tables for Inference: Model Checking by Enumeration

B <sub>1,1</sub>	B <sub>2,1</sub>	P <sub>1,1</sub>	P <sub>1,2</sub>	P <sub>2,1</sub>	P <sub>2,2</sub>	P <sub>3,1</sub>	R <sub>1</sub>		R <sub>2</sub>	R <sub>3</sub>	R <sub>4</sub>	R <sub>5</sub>	КВ
false	tru	ıe	True	true	true	false	False						
false	false	False	False	False	False	True	Tru	ıe	True	False	True	False	False
 false	 true	 false	 false	 false	 false	 false	 tru	ıe	 true	 false	 true	 true	 false
False	true	False	False	False	False	True	Tru	ıe	True	True	True	True	True
False	True	False	False	False	True	False	Tru	ue	True	True	True	True	True
False	true	False	false	false	true	true	tru	ıe	true	true	true	true	true
False	True	False	False	True	False	False	Tru	ле	False	False	True	True	False
true	fal	se	true	true	false	true	false						

Enumerate rows (different assignments to symbols); rows are possible models If KB is true in a row/model, check that  $\alpha$  is true, if not, entailment does not hold If entailment not broken over all rows where KB is true, then  $\alpha$  is entailed



# Inference by Enumeration

```
function TT-Entails?(KB, \alpha) returns true/false Inputs: KB (the knowledge base), a sentence in proposition logic \alpha, the query, a sentence in propositional logic symbols \leftarrow a list of the propositional symbols in KB and \alpha Return TT-Check-All(KB, \alpha, symbols,[])
```

```
function TT-Check-All(KB, α, symbols, []) returns true/false
   if Empty?(symbols) then
      if PL-TRUE?(KB, model) then return PL-TRUE(α, model)
      else return true;
   else do
      P ← First(symbols); rest ←Rest(symbols)
   return TT-Check-All(KB, α, rest, Extend(P, true, model)) and
      TT-Check-All(KB, α, rest, Extend(P, false, model))
```

 $O(2^n)$  for *n* symbols, problem is co-NP-complete.



#### **Proof Methods**

Proof methods divide into (roughly) two kinds:

#### Model checking:

- Truth table enumeration (always exponential in n)
- Improved backtracking, e.g., Davis-Putnam-Logemann-Loveland
- Backtracking with constraint propagation, backjumping
- Heuristic search in model space (sound but incomplete) e.g., min-conflicts, etc.

#### **Theorem Proving/Deductive Systems: Application of inference rules**

- Legitimate (sound) generation of new sentences from old
- Proof = a sequence of inference rule applications
- Typically requires translation of sentence into a normal form



# End of first part of Chap 7



# Logical Equivalence

Two sentences are logically equivalent iff true in same models:  $\alpha \equiv \beta$  if and only if  $a \models \beta$  and  $\beta \models \alpha$ 

$$(\alpha \wedge \beta)$$

Proof = a sequence of inference rule applications

Typically requires translation of sentence into a normal form