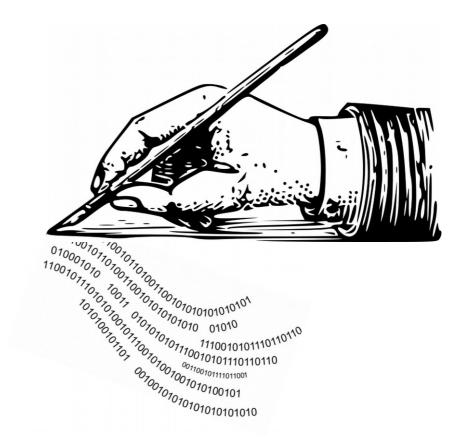
## CS 470 Spring 2018

Mike Lam, Professor



### Parallel Algorithm Development

(Foster's Methodology)

### Parallel program development

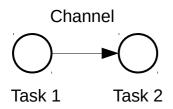
- Writing efficient parallel code is hard
- We've covered two generic paradigms ...
  - Shared-memory
  - Distributed message-passing
- ... and three specific technologies
  - Pthreads
  - OpenMP
  - MPI
- Given a problem, how do we approach the development of a parallel program that solves it?

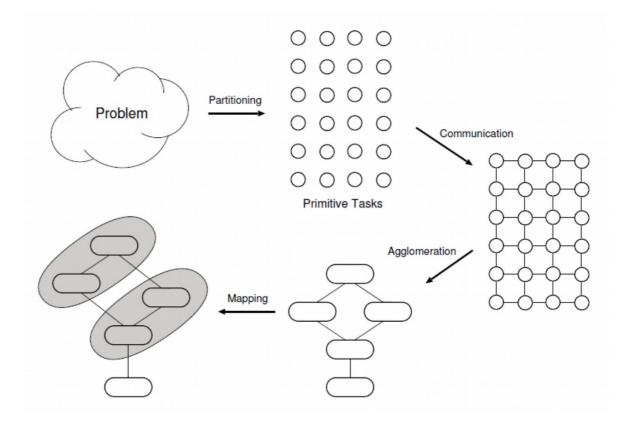
### Method vs. methodology

- Method: a systematic process or way of doing a task
- Methodology: analysis of methods relevant to a discipline
  - Literally: "the study of methods"
  - Goal: guidelines or best practices for a class of methods
- Parallel algorithms
  - There is no single **method** for creating efficient parallel algorithms
  - However, there are some good **methodologies** that can guide us
  - We will study one: Foster's methodology

### Foster's methodology

- Task: executable unit along with local memory and I/O ports
- Channel: message queue connecting tasks' input and output ports
- Drawn as a graph, tasks are vertices and channels are edges
- Steps:
  - 1) Partitioning
  - 2) Communication
  - 3) Agglomeration
  - 4) Mapping



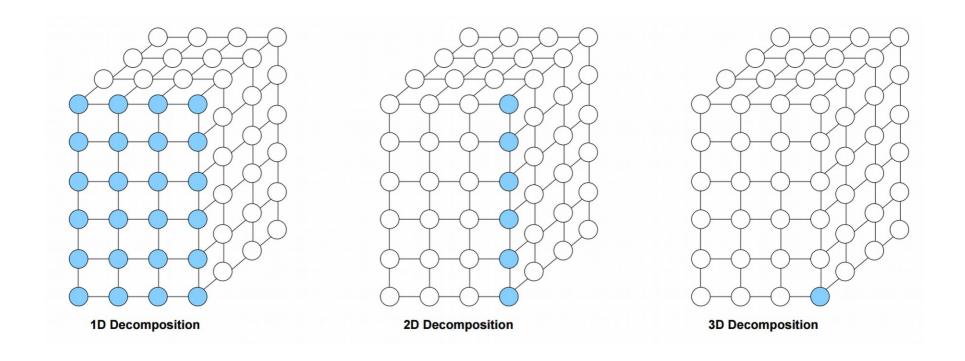


### Partitioning

- Goal: discover as much parallelism as possible
- Divide computation into as many primitive tasks as possible
  - Avoid redundant computation
  - Primitive tasks should be roughly the same size
  - Number of tasks should increase as the problem size increases
    - This helps ensure good scaling behavior

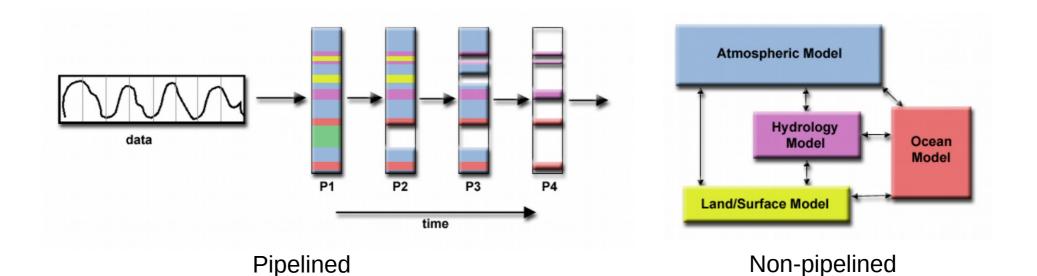
### Partitioning

- Domain ("data") decomposition
  - Break tasks into segments of various granularities by data



### Partitioning

- Functional ("task") decomposition
  - Separation by task type
  - Domain/data decomposition can often be used inside of individual tasks

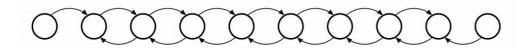


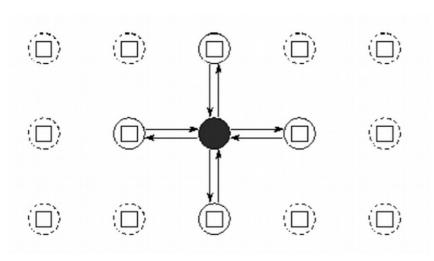
### Communication

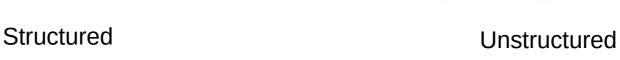
- Goal: minimize overhead
- Identify which tasks must communicate and how
  - Local (few tasks) vs. global (many tasks)
  - Structured (regular) vs. unstructured (irregular)
  - Prefer local, structured communication
  - Tasks should perform similar amounts of communication
    - This helps with load balancing
  - Communication should be concurrent wherever possible

### Communication

• Examples of local communication:

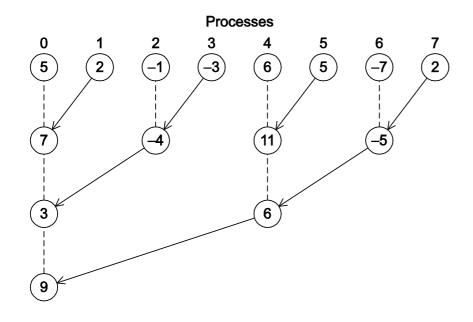


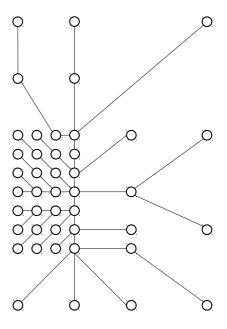




### Communication

• Examples of global communication:





Structured

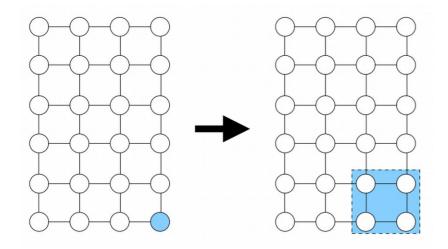
Unstructured

### Agglomeration

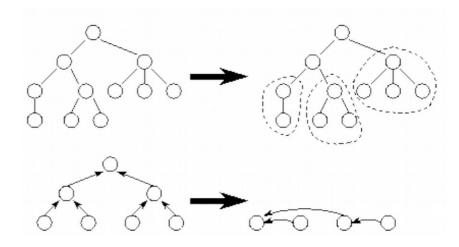
- Goal: Reduce messages and simplify programming
- Combine tasks into groups, increasing locality
  - Groups should have similar computation and communication costs
  - Task counts should still scale with processor count and / or problem size
  - Minimize software engineering costs
    - Agglomeration can prevent code reuse

## Agglomeration

### • Examples:



Agglomeration of four local tasks



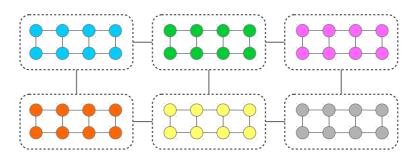
Agglomeration of tree-based tasks

### Mapping

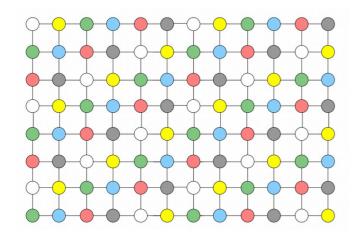
- Goal: minimize execution time
  - Alternately: maximize processor utilization
  - On a distributed system: minimize communication
- Assign tasks (or task groups) to processors/nodes
  - Block vs. cyclic
  - Static vs. dynamic
- Strategies:
  - 1) Place concurrent tasks on different nodes
  - 2) Place frequently-communicating tasks on the same node
- Problem: these strategies are often in conflict!
  - The general problem of optimal mapping is NP-complete

# Mapping

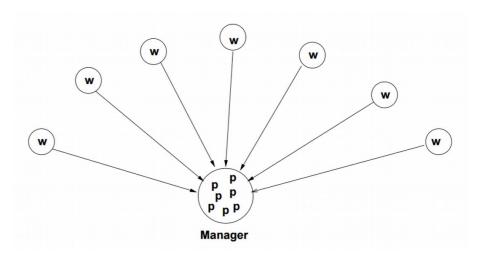
• Examples:



**Block mapping** 



Cyclic mapping



Dynamic mapping

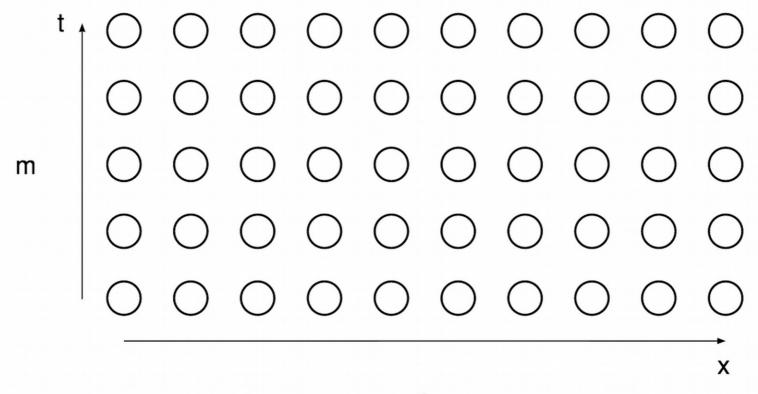
#### Problem

- General statement: Determine the temperature changes in a thin cylinder of uniform material with constant-temperature boundary caps over a given time period, given the size of the cylinder and its initial temperature
- General solution: solve partial differential equation(s)
  - Often too difficult or expensive to solve analytically
- Approximate solution: finite difference method
  - Discretize space (1d grid) and time (ms)
- Goal: Parallelize this solution, using Foster's methodology as a guide

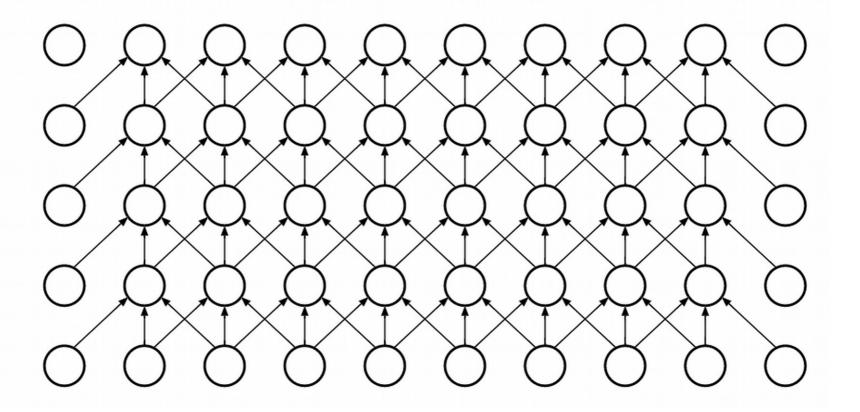
#### Partitioning:

Make each T(x, t) computation a primitive task.

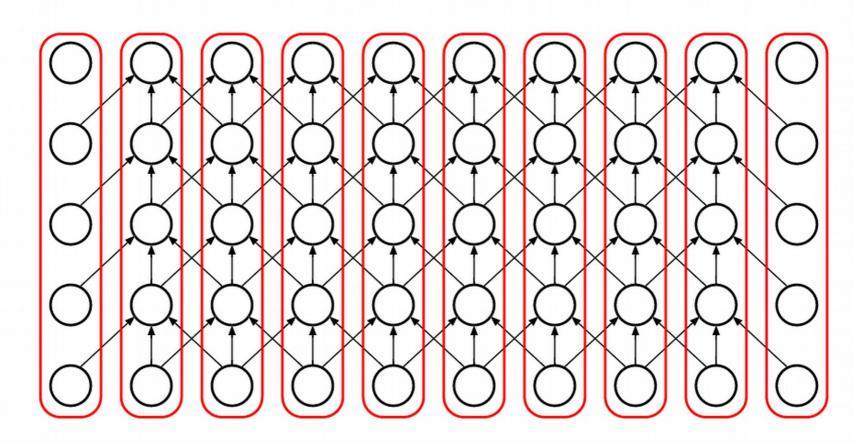
⇒ 2-dimensional domain decomposition



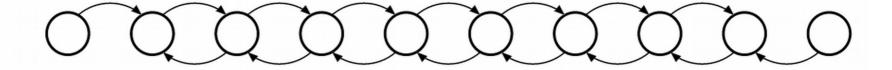
#### Communication:



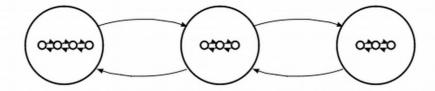
### Agglomeration:



### Agglomeration:



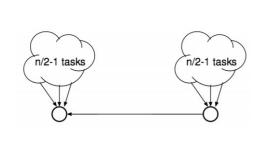
### Mapping:

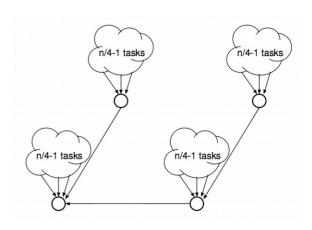


- Problem: Determine the maximum value among some large set of given values
  - Special case of a reduction
- Goal: Parallelize this solution, using Foster's methodology as a guide

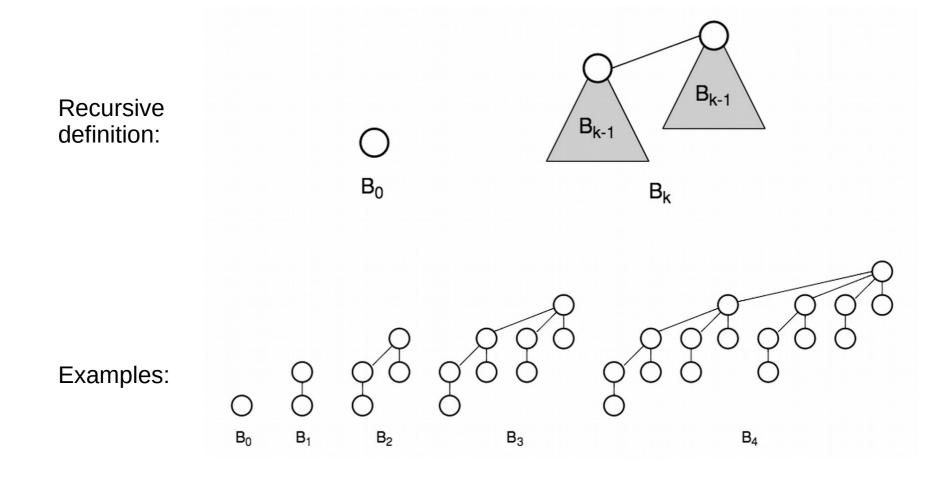
- Partitioning: each value is a primitive task
  - (1d domain decomposition)
  - One task (root) will compute final solution
- Communication: divide-and-conquer
  - Root task needs to compute max after n-1 tasks
  - Keep splitting the input space in half





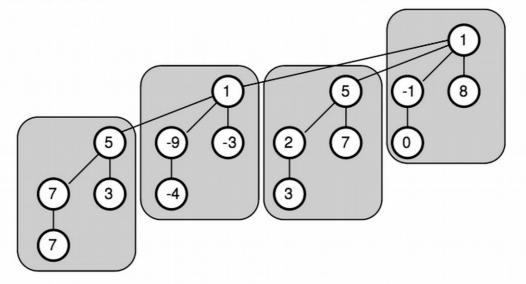


- Binomial tree with  $n = 2^k$  nodes
  - (remember merge sort in P2?)



#### Agglomeration:

Group *n* leafs of the tree:



#### Mapping:

The same (actually, in the agglomeration phase, use n such that you end up with p tasks).

### Random number generation

- Goal: Generate psuedo-random numbers in a distributed way
- Problem: We wish to retain some notion of reproducibility
  - In other words: results should be deterministic, given the RNG seed
  - This means we can't depend on the ordering of distributed communications
- Problem: We wish to avoid duplicated series of generated numbers
  - This means we can't just use the same generator in all processes



### Random number generation

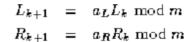
- Naive solution:
  - Generate all numbers on one node and scatter them (a la P2)
  - Too slow!
- Can we do better? (Foster's)
  - Generating each random number is a task
  - Channels between subsequent numbers from the same seed
  - Tweak communication & agglomeration
  - Minimize dependencies

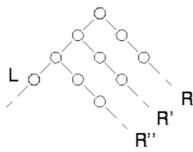


### Random number generation

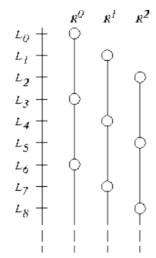
#### Goal:

Uniform randomness and reproducibility





**Figure 10.1:** The random tree method. Two generators are used to construct a tree of random numbers. The right generator is applied to elements of the sequence L generated by the left generator to generate new sequences R, R', R'', etc.



**Figure 10.2:** The leapfrog method with n=3. Each of the three right generators selects a disjoint subsequence of the sequence constructed by the left generator's sequence.

### Common paradigms

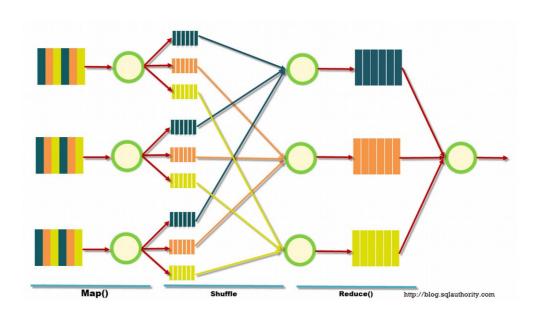
- Grid/mesh-based nearest-neighbor simulation
  - Often includes math-heavy computations
    - Linear algebra and systems of equations
    - Dense vs. sparse matrices
  - Newer: adaptive mesh and multigrid simulations
- Worker pools / task queues
  - Newer: adaptive cloud computing
- Pipelined task phases
  - Newer: MapReduce
- Divide-and-conquer tree-based computation
  - Often combined with other paradigms (worker pools and pipelines)

### MapReduce

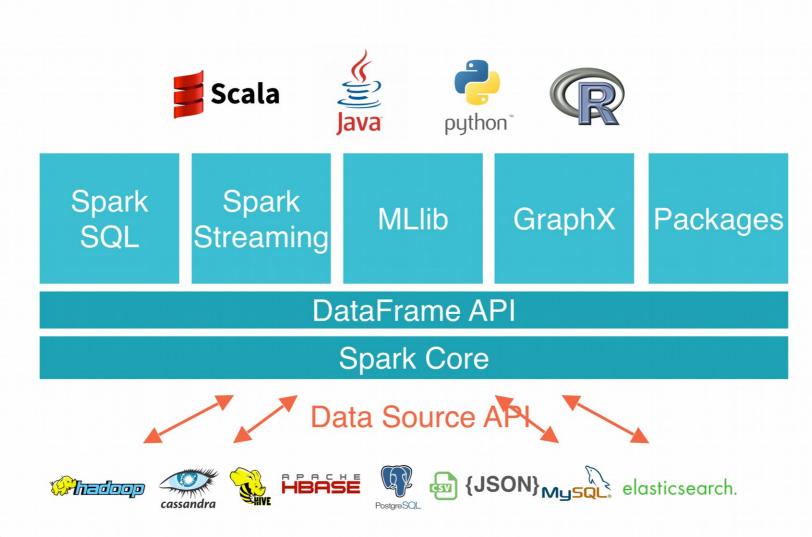
- Parallel/distributed system paradigm for "big data" processing
  - Uses a specialized file system
  - Originally developed at Google (along with GFS)
  - Currently popular: Apache Hadoop and HDFS
    - General languages: Java, Python, Ruby, etc.
    - Specialized languages: Pig (data flow language) or Hive (SQL-like)
    - Growing quickly: Apache Spark (more generic w/ in-memory processing)

#### Phases

- Map (process local data)
- Shuffle (distributed sort)
- Reduce (combine results)



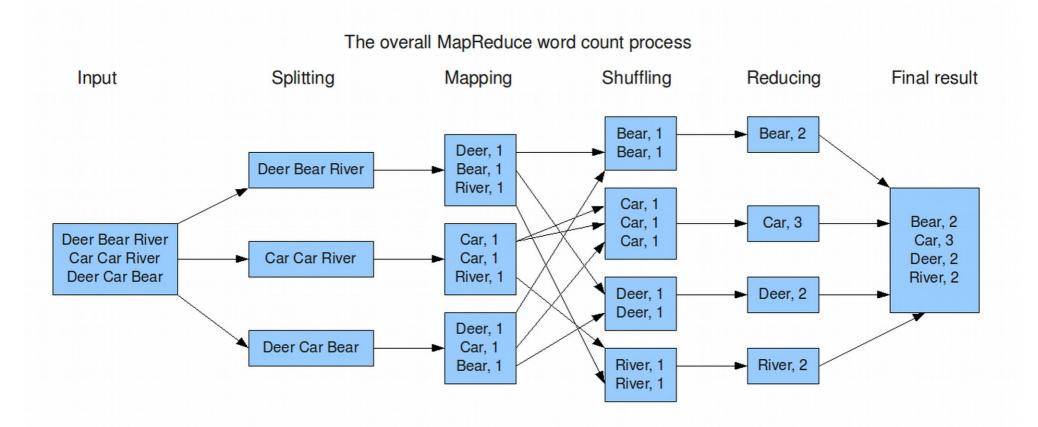
## Apache Spark





### MapReduce

Word count example



## Apache Hadoop (Java)

```
public class WordCount {
public static class Map extends MapReduceBase implements Mapper<LongWritable, Text, Text, IntWritable> {
   private final static IntWritable one = new IntWritable(1);
   private Text word = new Text();
   public void map(LongWritable key, Text value, OutputCollector<Text,</pre>
                   IntWritable> output, Reporter reporter) throws IOException {
    String line = value.toString();
    StringTokenizer tokenizer = new StringTokenizer(line);
    while (tokenizer.hasMoreTokens()) {
      word.set(tokenizer.nextToken());
      output.collect(word, one);
  }
}
public static class Reduce extends MapReduceBase implements Reducer<Text, IntWritable, Text, IntWritable> {
  public void reduce(Text key, Iterator<IntWritable> values,
                      OutputCollector<Text, IntWritable> output, Reporter orter) throws IOException {
     int sum = 0;
    while (values.hasNext()) {
      sum += values.next().get();
    output.collect(key, new IntWritable(sum));
```

### Apache Spark (Python)

#### **WORD COUNT**

#### MONTE CARLO PI