

CS 261

Fall 2016

Mike Lam, Professor

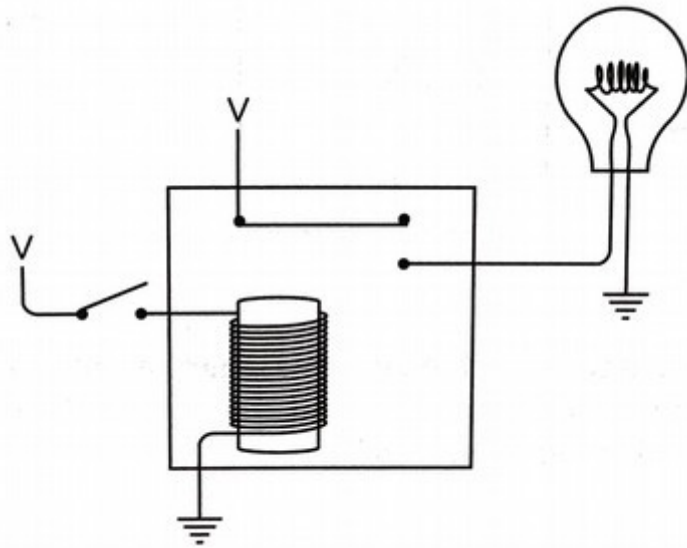
Logic Gates

The final frontier

- Java programs running on Java VM
- C programs compiled on Linux
- Assembly / machine code on CPU + memory
- ???
- Switches and electric signals

Aside: Relays

- From “Code” recommended reading:

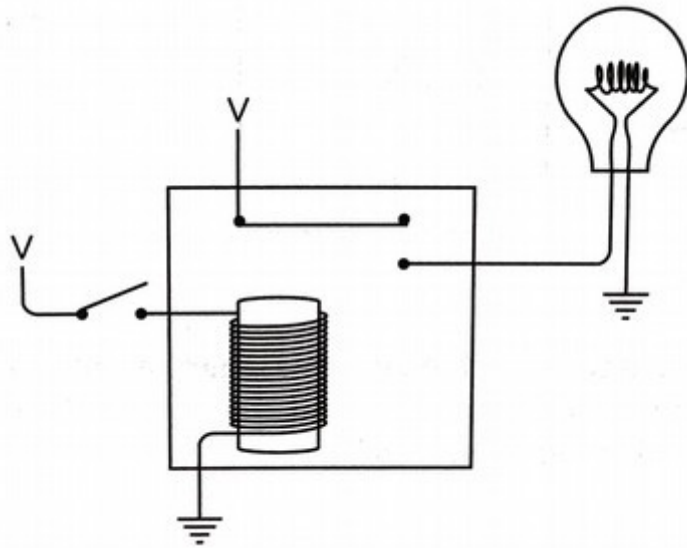


Relay

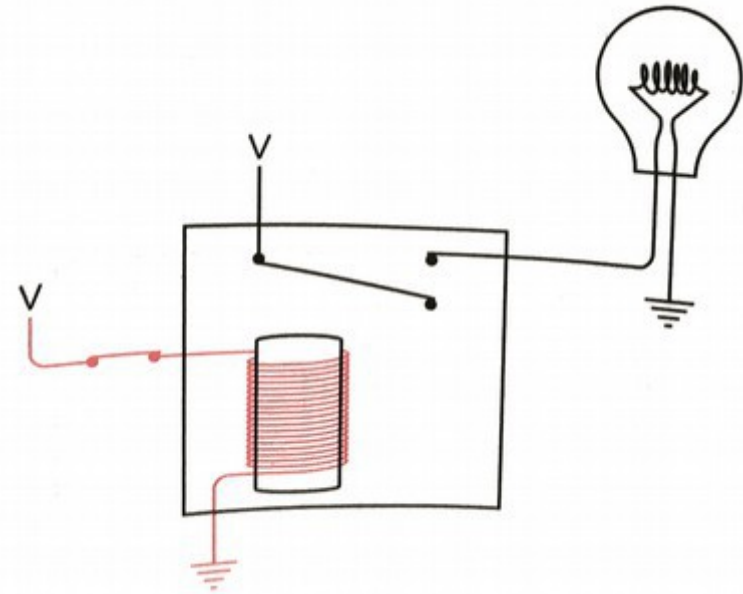
Question: what happens if we connect the light bulb to the other contact?

Aside: Relays

- From “Code” recommended reading:



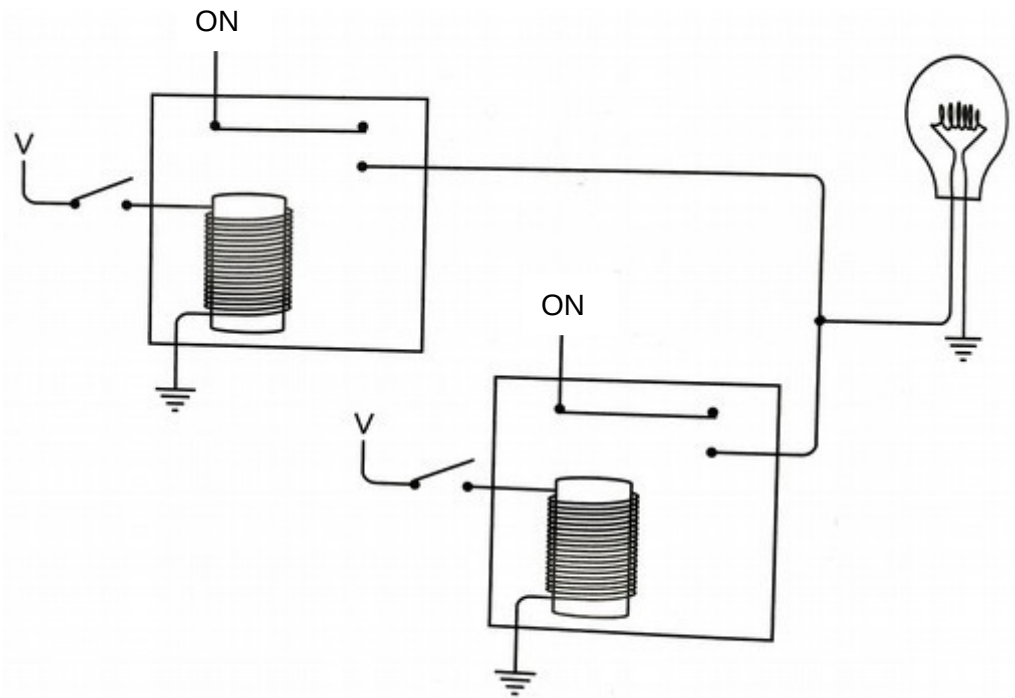
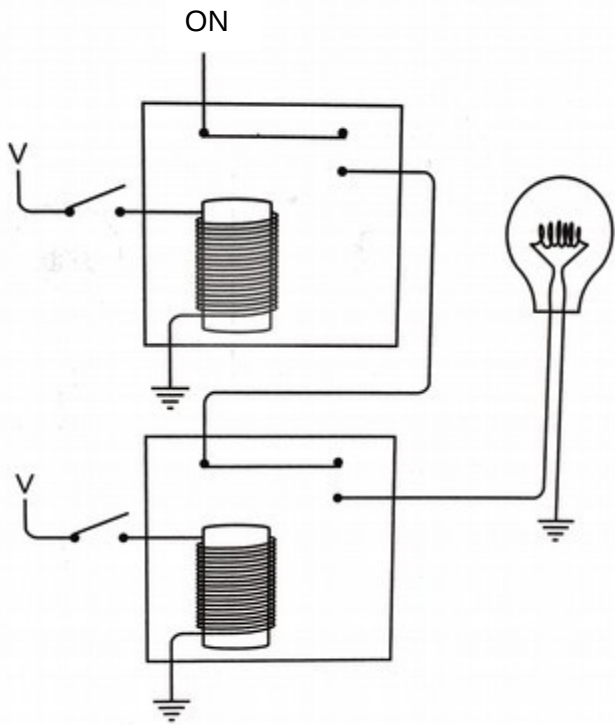
Regular relay



Inverted relay (NOT)

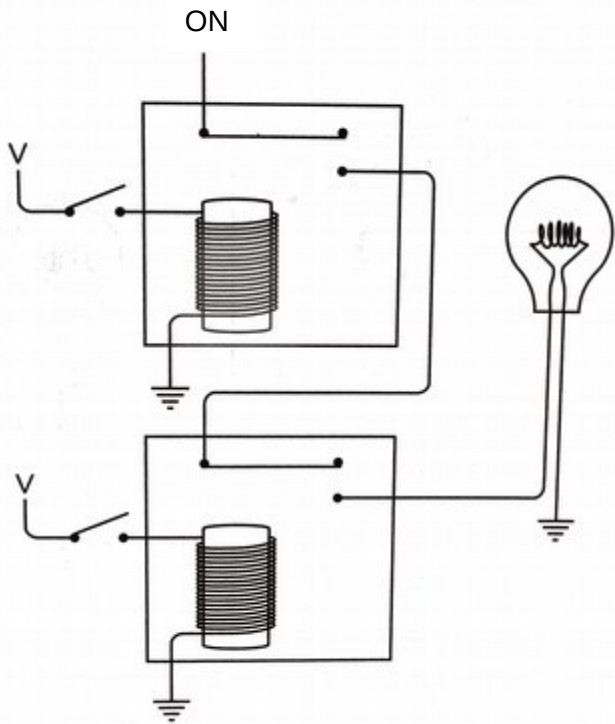
Aside: Relays

- From “Code” recommended reading:

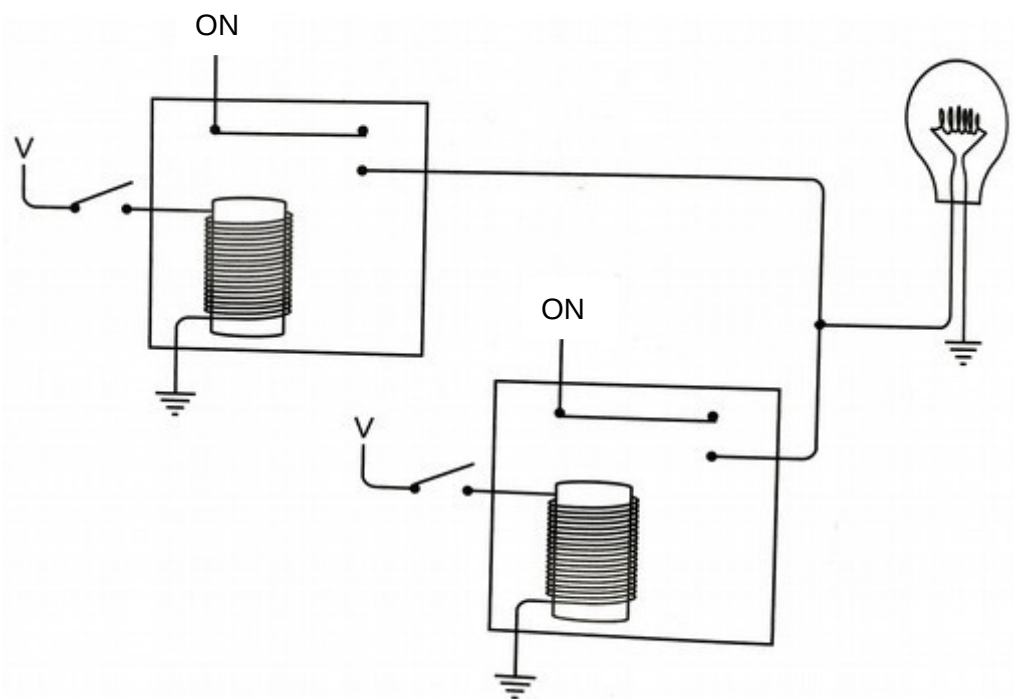


Aside: Relays

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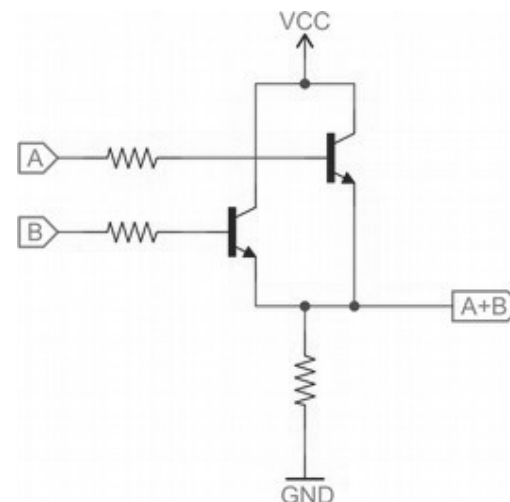
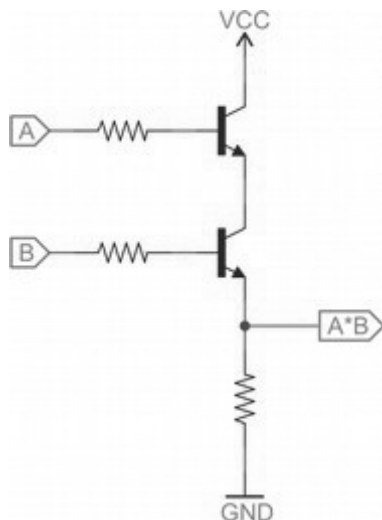
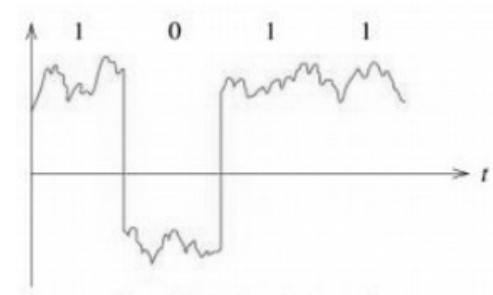
Relays in series (AND)



Relays in parallel (OR)

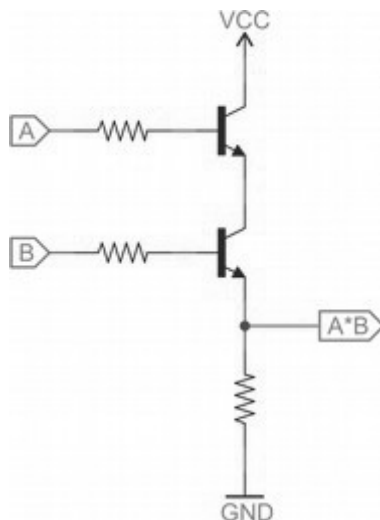
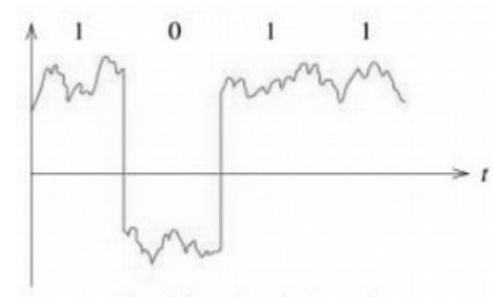
Digital hardware

- Digital signals are transmitted via electric signals by varying voltages
 - 1.0 V (high) = binary 1
 - 0.0 V (low) = binary 0
 - Use a threshold to distinguish

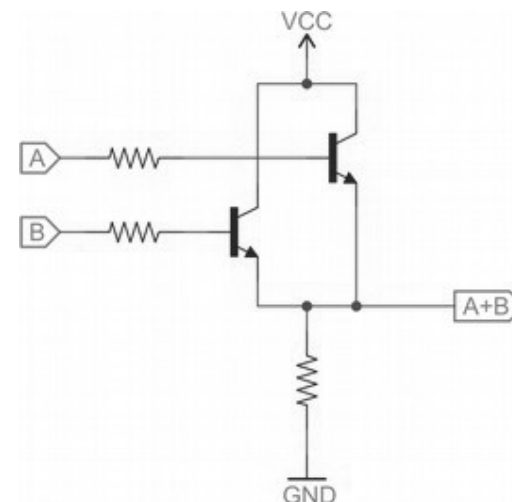


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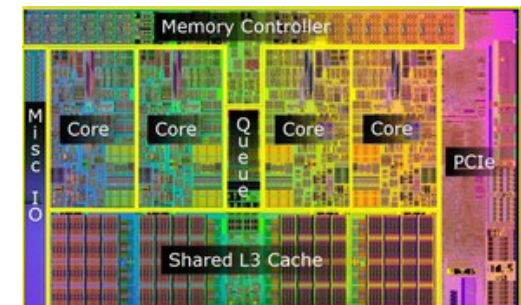
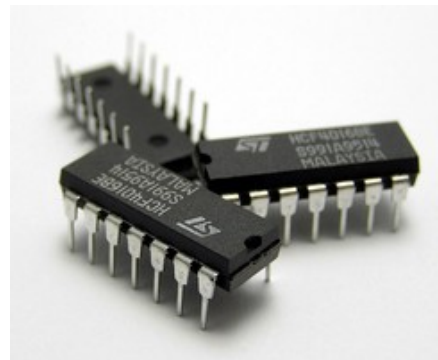
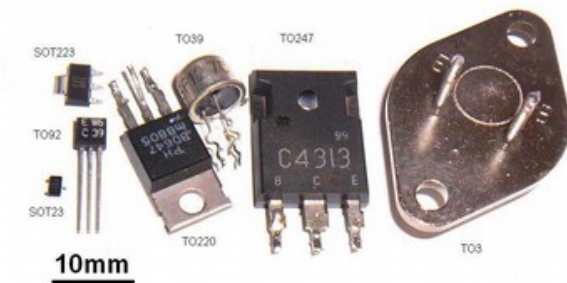
AND



OR

Transistors

- **Transistors** are the fundamental hardware component of computing
 - Similar to relays; replaced vacuum tubes
 - Smaller, more reliable, and use less energy
 - Primary functions: switching and amplification
 - Mostly silicon-based semiconductors now
 - **Metal–Oxide–Semiconductor Field-Effect Transistor** (MOSFET)
 - n-channel (“on” when $V_{\text{gate}} = 1\text{V}$) vs. p-channel (“off” when $V_{\text{gate}} = 1\text{V}$)
 - Mass-produced on **integrated circuit** chips
 - For convenience, we abstract their behavior using **logic gates**



Logic gates

- Primary gates:



&	0	1
0	0	0
1	0	1



	0	1
0	0	1
1	1	1



!	
0	1
1	0



	0	1
0	1	1
1	1	0



	0	1
0	1	0
1	0	0



^	0	1
0	0	1
1	1	0

Logic gates

- Primary gates:



&	0	1
0	0	0
1	0	1

AND



	0	1
0	0	1
1	1	1

OR



!	
0	1
1	0

NOT



	0	1
0	1	1
1	1	0

NAND



	0	1
0	1	0
1	0	0

NOR



^	0	1
0	0	1
1	1	0

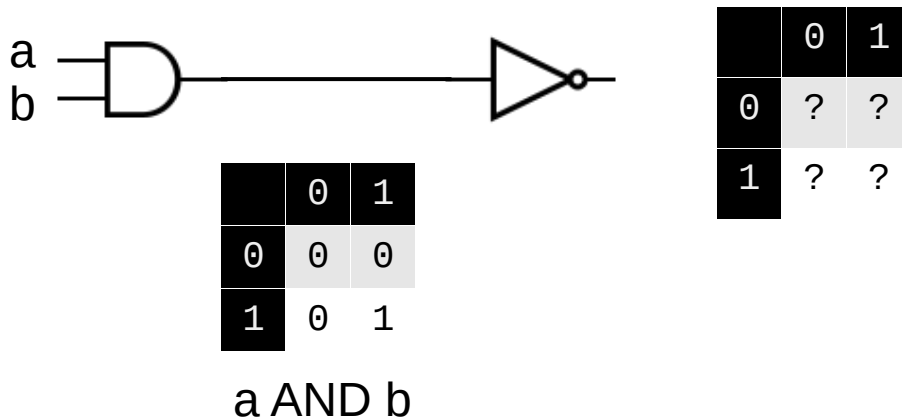
XOR

Basic combinatorial circuits

- **Circuits** are formed by connecting gates together
 - Textbook uses Hardware Description Language (HDL)
 - Equivalent to **boolean formulas** or **functions**
 - $f(g(x, y))$ means apply “operation f to the result of operation g on x and y ”
 - In a diagram: $x, y \rightarrow g \rightarrow f$ (i.e., ordering is g first, then f)

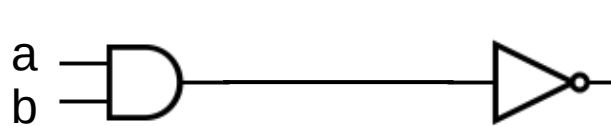
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 - $f(g(x, y))$ means apply “operation f to the result of operation g on x and y ”
 - In a diagram: $x, y \rightarrow g \rightarrow f$ (i.e., ordering is g first, then f)
 - NAND example: (similarly for NOR)
 - Infix/boolean notation: $a \text{ NAND } b = !(a \& b)$
 - Function notation: $\text{NAND}(a, b) = \text{NOT}(\text{AND}(a, b))$



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	0	1
0	0	0
1	0	1

a AND b

	0	1
0	1	1
1	1	0

NOT (a AND b)



	0	1
0	1	1
1	1	0

a NAND b

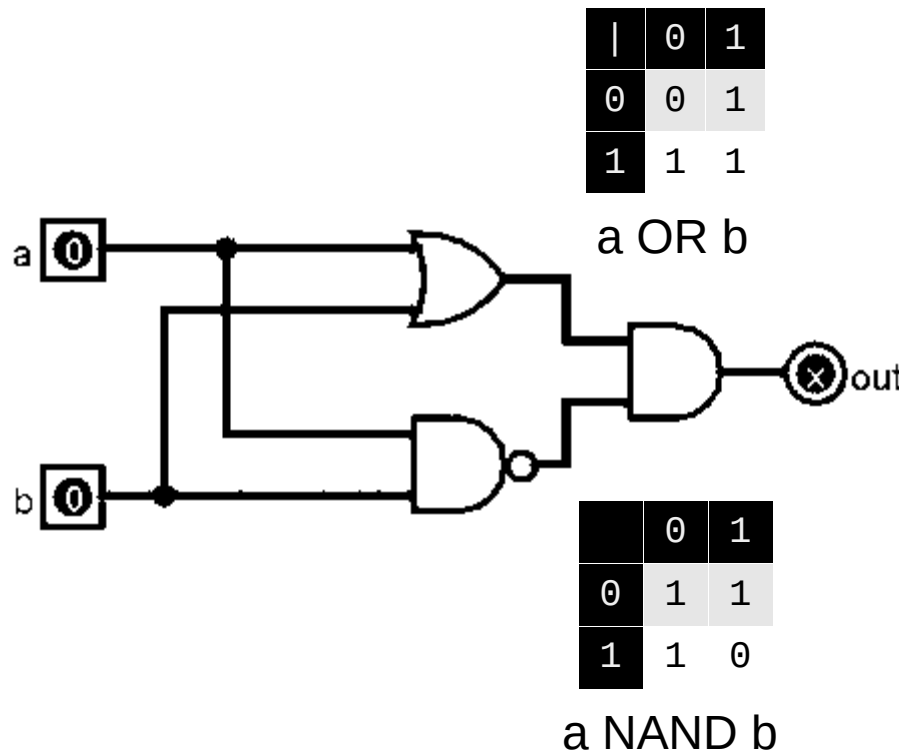
Basic combinatorial circuits

- Circuits are **equivalent** if the truth tables are the same
 - $a \text{ XOR } b = (a \text{ OR } b) \text{ AND } (a \text{ NAND } b)$
 - $\text{XOR}(a, b) = \text{AND}(\text{OR}(a,b), \text{NAND}(a,b))$



\wedge	0	1
0	0	1
1	1	0

XOR



	0	1
0	0	1
1	1	1

	0	1
0	1	1
1	1	0

	0	1
0	?	?
1	?	?

(a OR b) AND
(a NAND b)

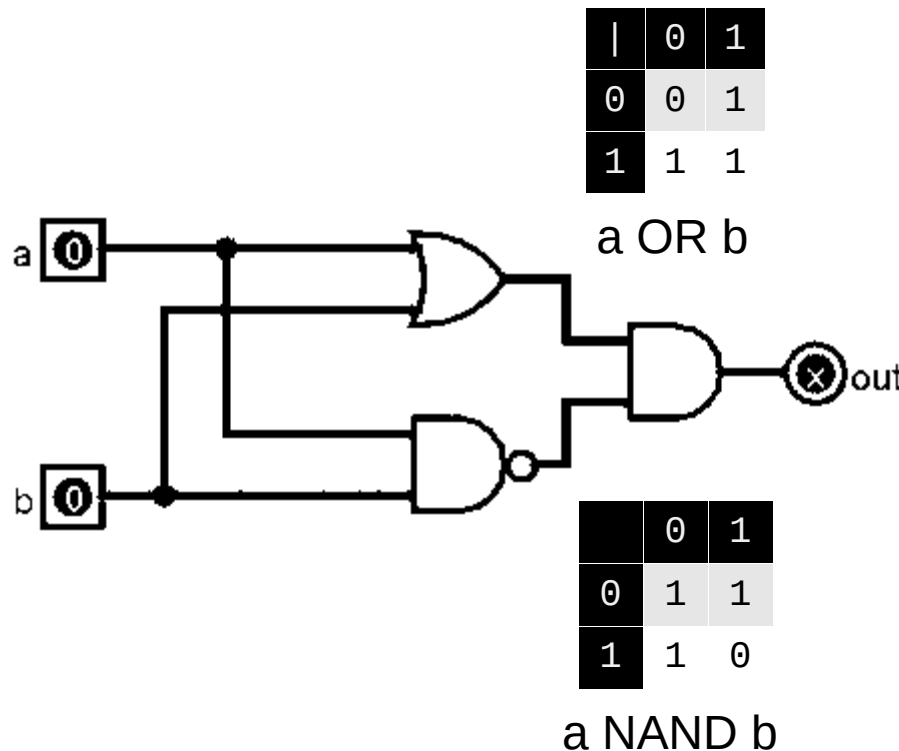
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\wedge	0	1
0	0	1
1	1	0

XOR



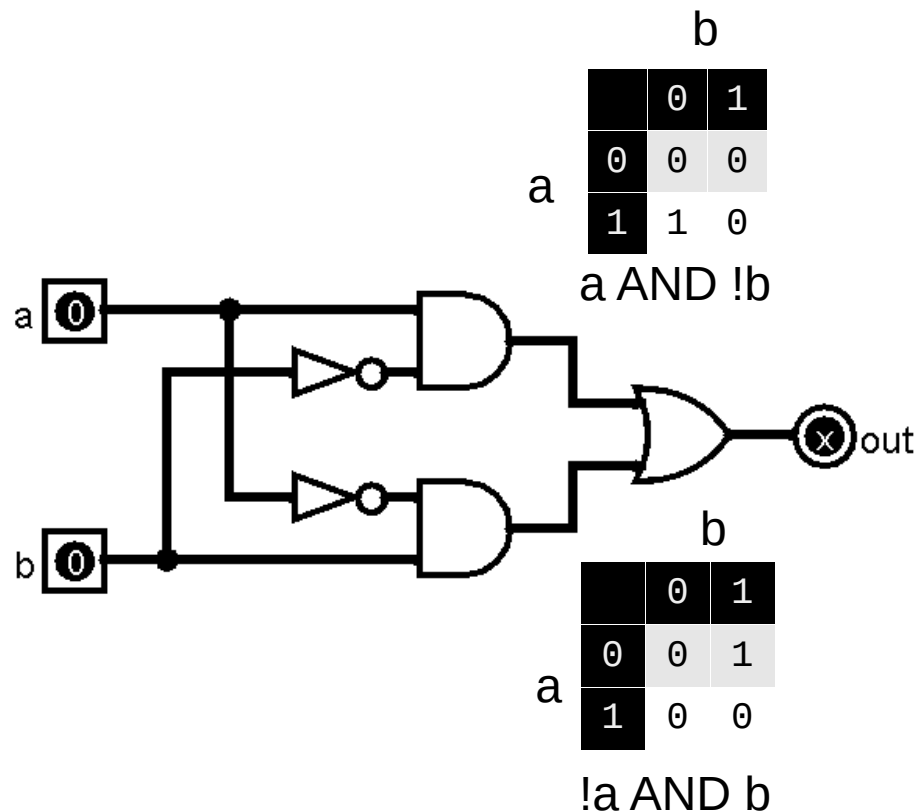
	0	1
0	0	1
1	1	1

	0	1
0	0	1
1	1	0

(a OR b) AND
(a NAND b)

Basic combinatorial circuits

- Circuits can be equivalent even if the structure is different
 - $f(a, b) = (a \text{ AND } !b) \text{ OR } (!a \text{ AND } b)$
 - What is this equivalent to?



	0	1
0	?	?
1	?	?

(a AND !b) OR
(!a AND b)

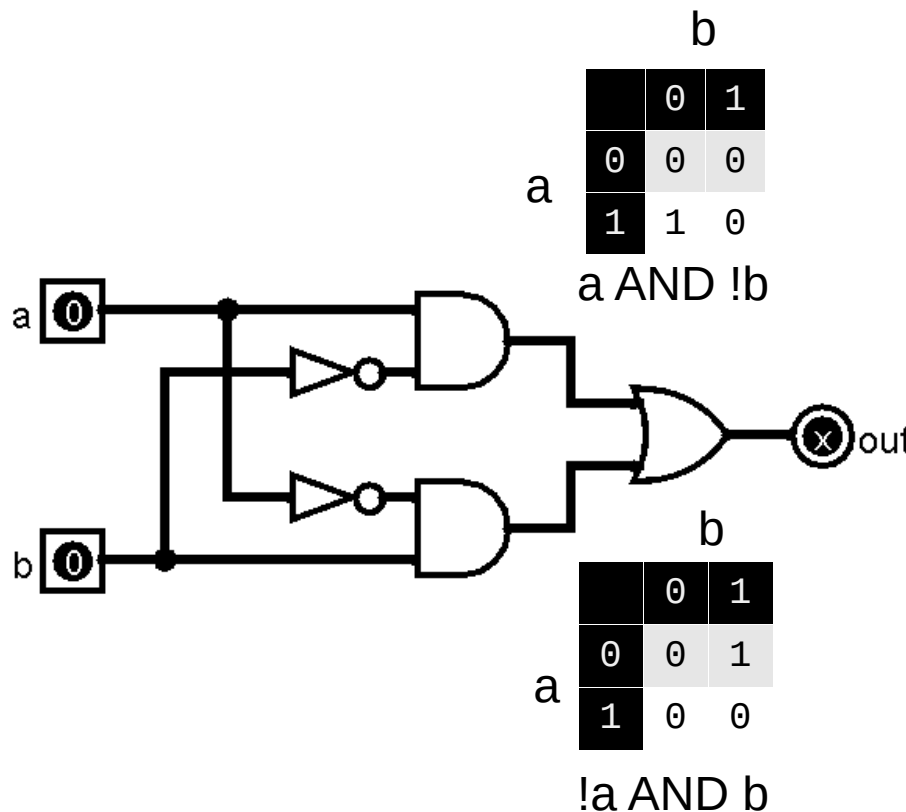
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 - $f(a, b) = (a \text{ AND } !b) \text{ OR } (!a \text{ AND } b)$
 - What is this equivalent to?



\wedge	0	1
0	0	1
1	1	0

XOR



	0	1
0	0	1
1	1	0

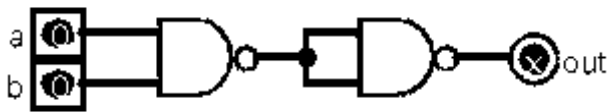
$(a \text{ AND } !b) \text{ OR } (!a \text{ AND } b)$

Important properties

- Identity: $a \text{ AND } 1 = a$ $(a \text{ OR } 0) = a$
 - Constants: $a \text{ AND } 0 = 0$ $(a \text{ OR } 1) = 1$
 - Also: $a \text{ NAND } 0 = 1$ $(a \text{ NOR } 1) = 0$
 - Inverses: $a \text{ NAND } 1 = !a$ $(a \text{ NOR } 0) = !a$
 - Also: $a \text{ NAND } a = !a$ $a \text{ NOR } a = !a$
 - Double inverse: $!!a = a$
 - Or: $\text{NOT}(\text{NOT}(a)) = a$
 - De Morgan's law: $!(a \ \& \ b) = !a \ | \ !b$
 - Alternatively: $!(a \ | \ b) = !a \ \& \ !b$
- (remember this from CS 227?)*

Universal gates

- NAND and NOR gates are **universal**
 - Each one alone can reproduce all other gates
 - Example: **a AND b** = $a \& b = \neg(\neg(a \& b)) = \neg(a \text{ NAND } b) = (a \text{ NAND } b) \text{ NAND } (a \text{ NAND } b)$



	0	1
0	1	1
1	1	0

a NAND b

	0	1
0	0	0
1	0	1

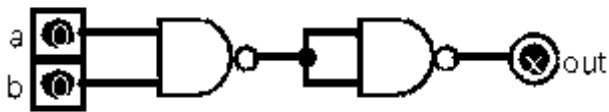
(a NAND b) NAND
(a NAND b)

	0	1
0	0	0
1	0	1

a AND b

Universal gates

- NAND and NOR gates are **universal**
 - Each one alone can reproduce all other gates
 - Example: **$a \text{ AND } b = a \& b = \neg(\neg(a \& b)) = \neg(a \text{ NAND } b) = (a \text{ NAND } b) \text{ NAND } (a \text{ NAND } b)$**
 - Similarly: **$a \text{ AND } b = \neg(\neg(a \& b)) = \neg(\neg a \mid \neg b) = \neg a \text{ NOR } \neg b = (a \text{ NOR } a) \text{ NOR } (b \text{ NOR } b)$**



	0	1
0	1	1
1	1	0

a NAND b

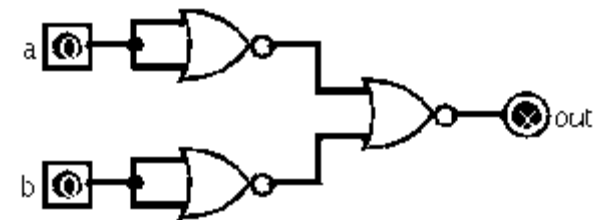
	0	1
0	0	0
1	0	1

(a NAND b) NAND
(a NAND b)



	0	1
0	0	0
1	0	1

a AND b

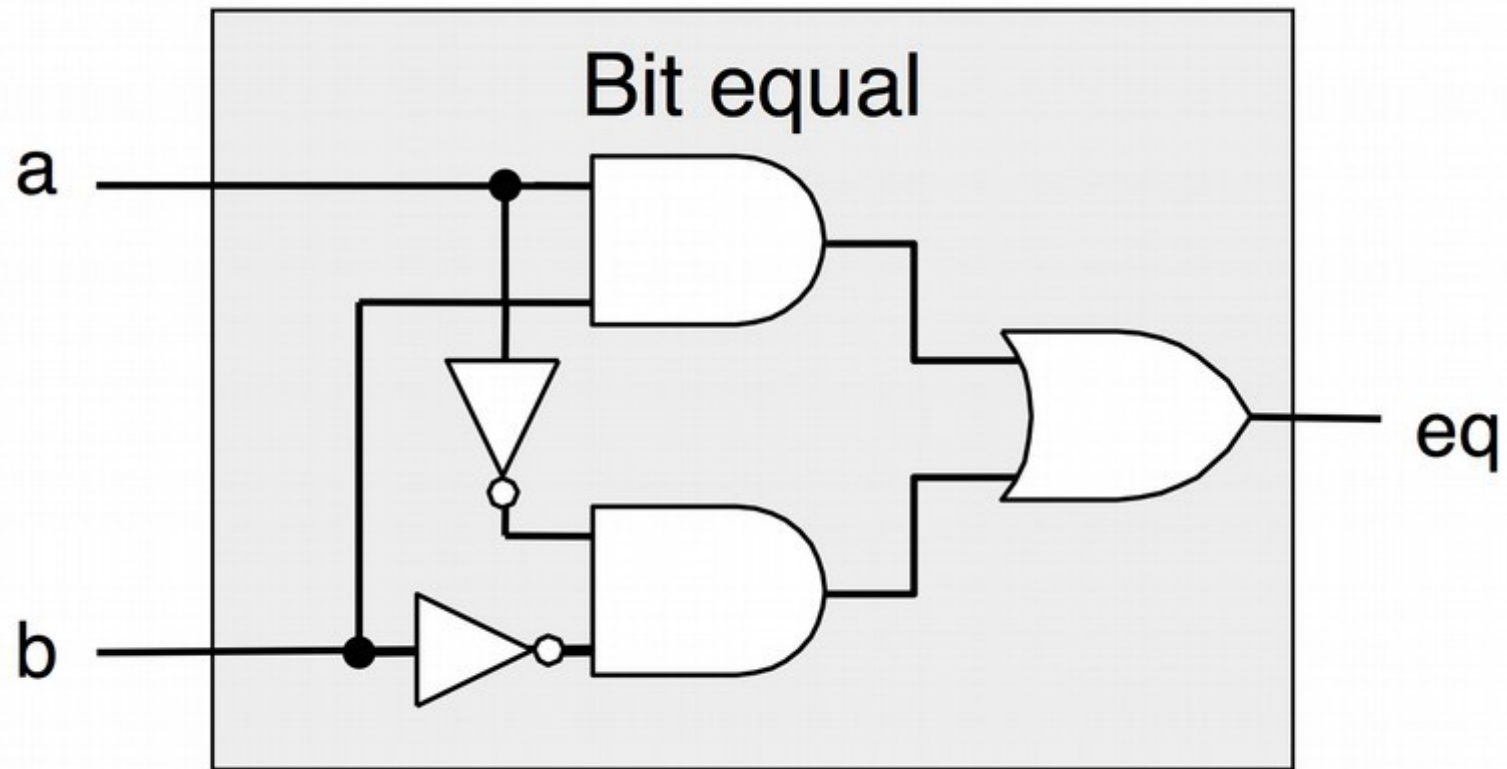


(a NOR a) NOR
(b NOR b)

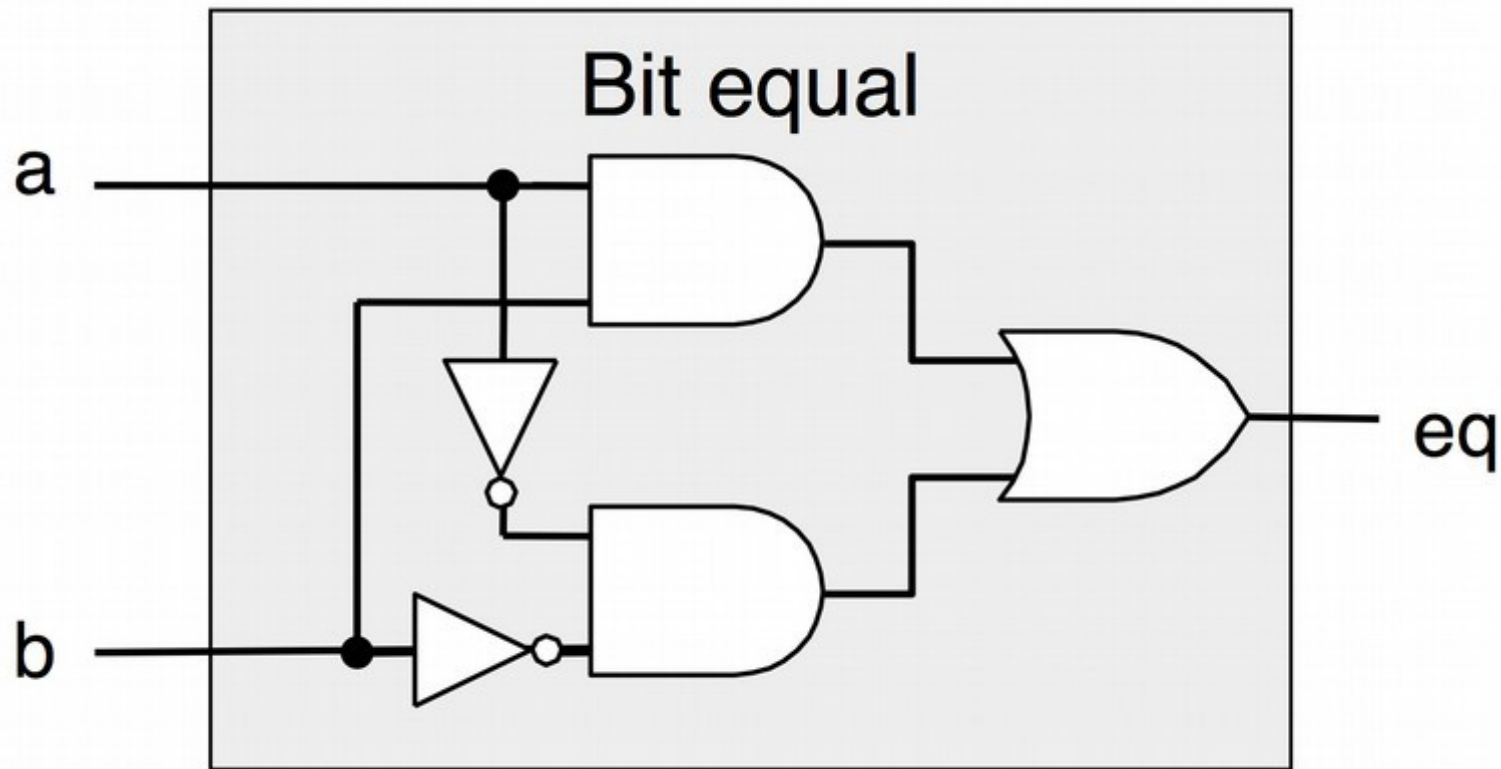
Computation

- Identify circuits that perform useful computation
 - Testing bits to see if they're equal
 - Selecting between multiple inputs
 - Adding or subtracting bits
 - Bitwise operations (AND, OR, XOR)
 - Make them work on bytes instead of bits

Equality



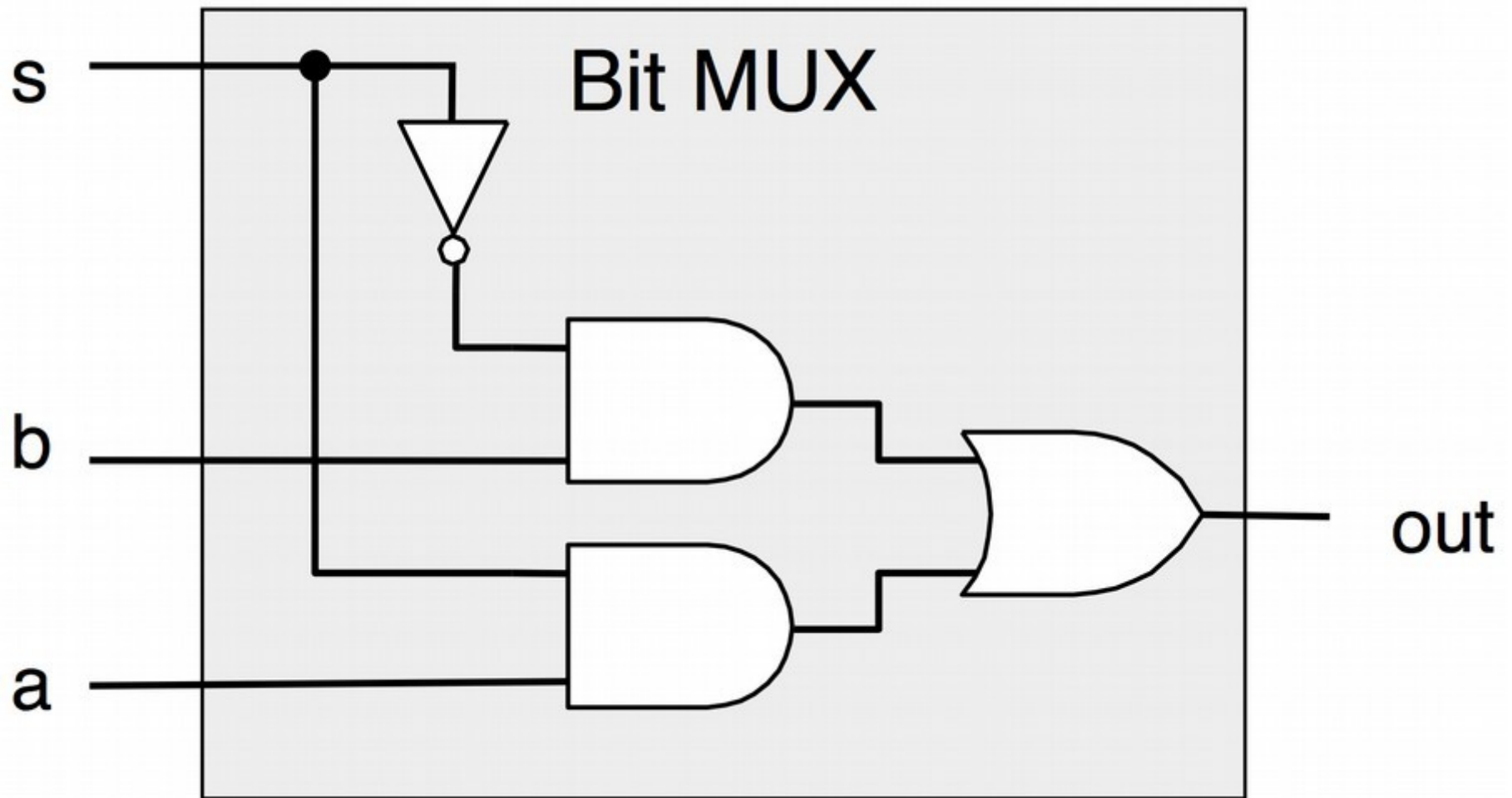
Equality



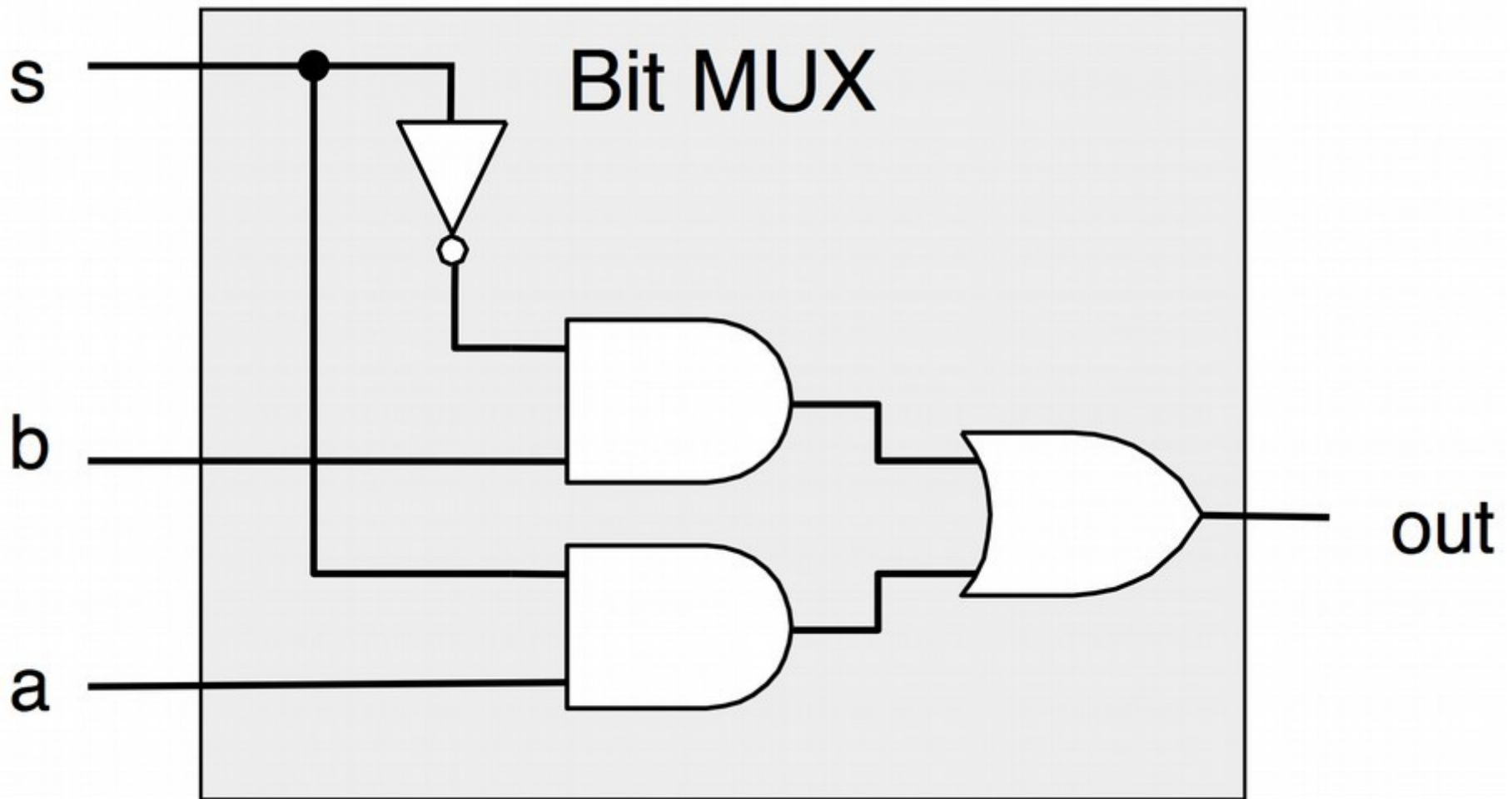
$$a \text{ EQ } b = (a \ \& \ b) \ | \ (!a \ \& \ !b)$$

$$\text{EQ}(a, b) = \text{OR}(\text{AND}(a, b), \text{AND}(\text{NOT}(a), \text{NOT}(b)))$$

Multiplexor (“selector”)



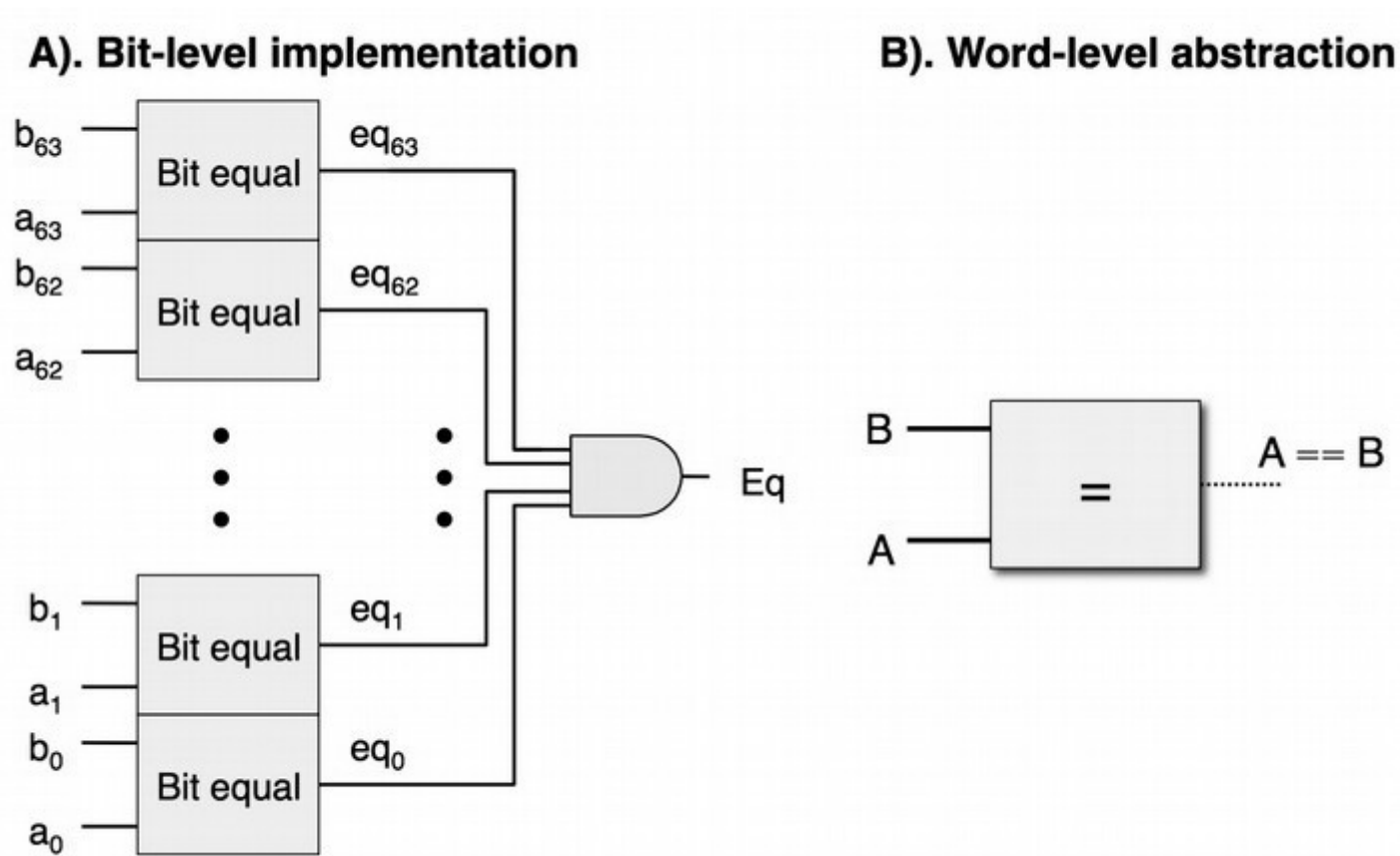
Multiplexor (“selector”)



$$\text{MUX}(a, b, s) = (s \& a) | (!s \& b)$$

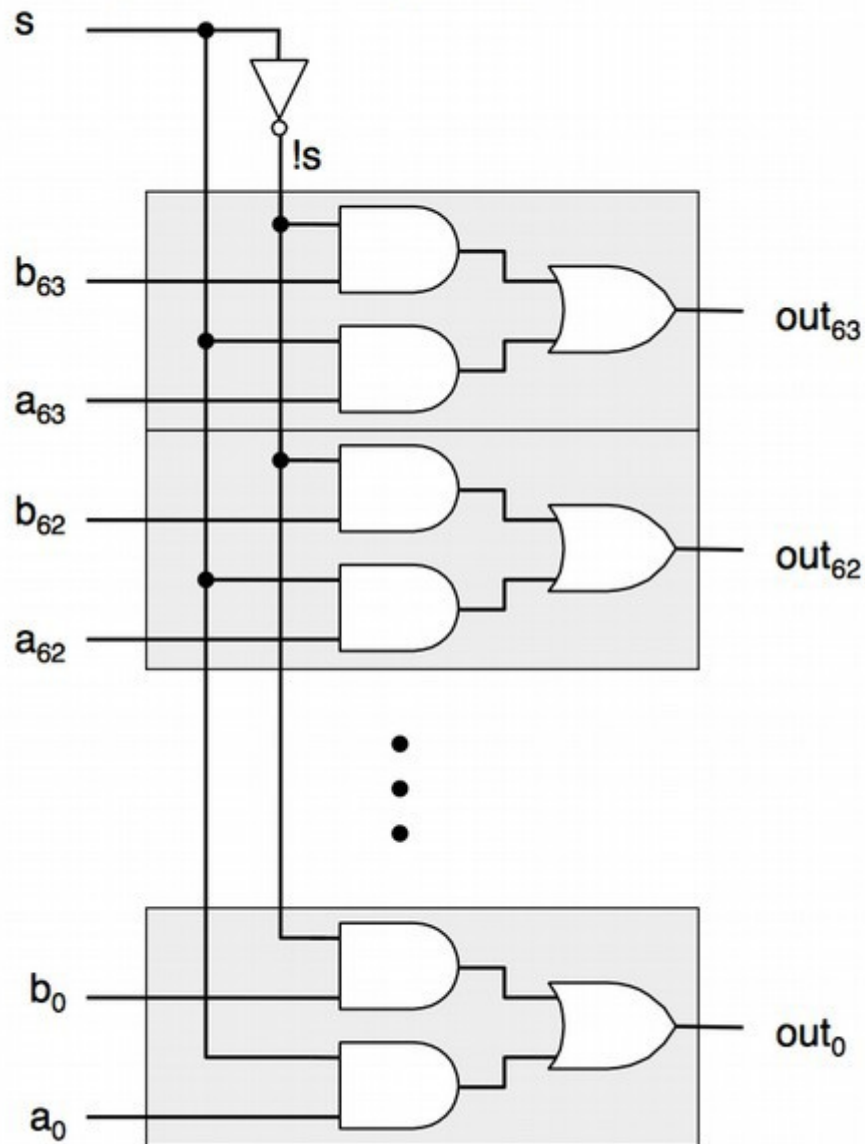
Abstraction

- Name circuits, then use them to build more complex circuits
 - E.g., use bit-level EQ to build a word-level equality circuit:

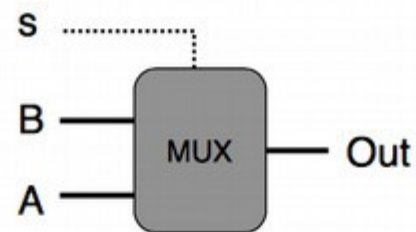


Word-level 2-way multiplexer

A). Bit-level implementation

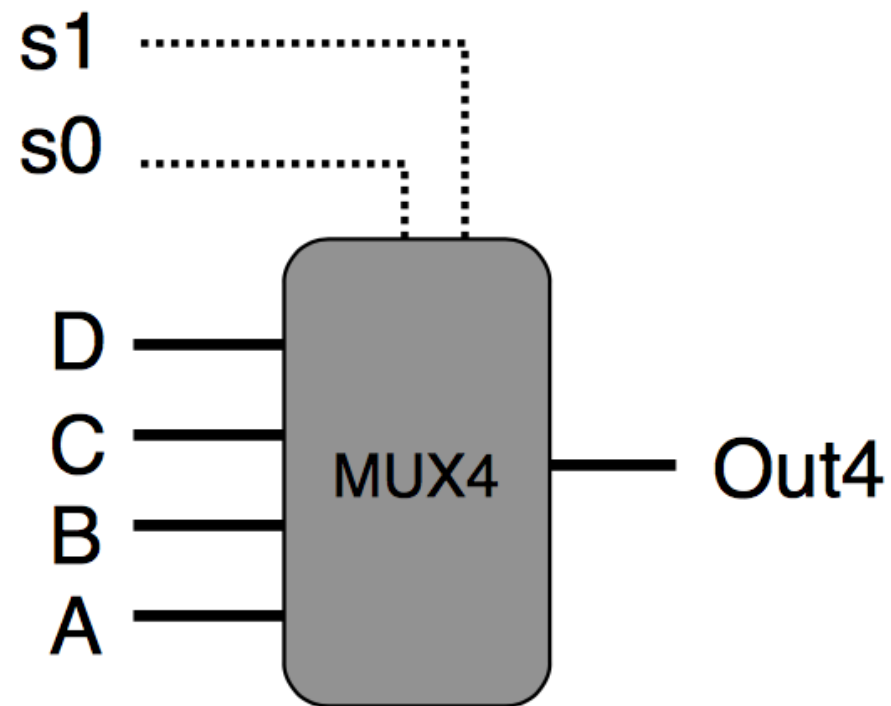


B). Word-level abstraction

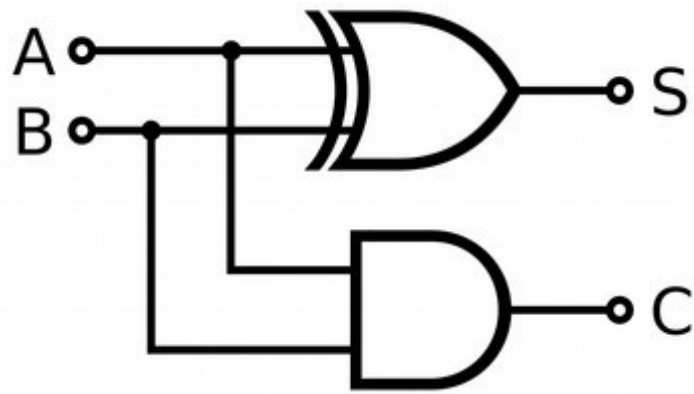


```
int Out = [  
    s : A;  
    1 : B;  
];
```

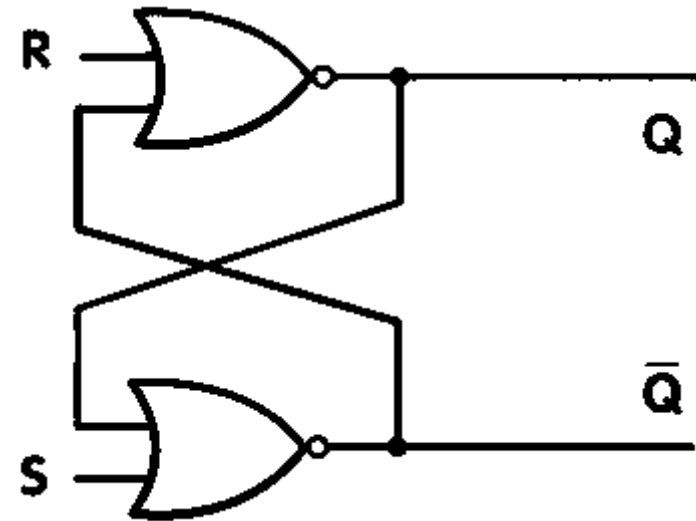
Word-level 4-way multiplexer



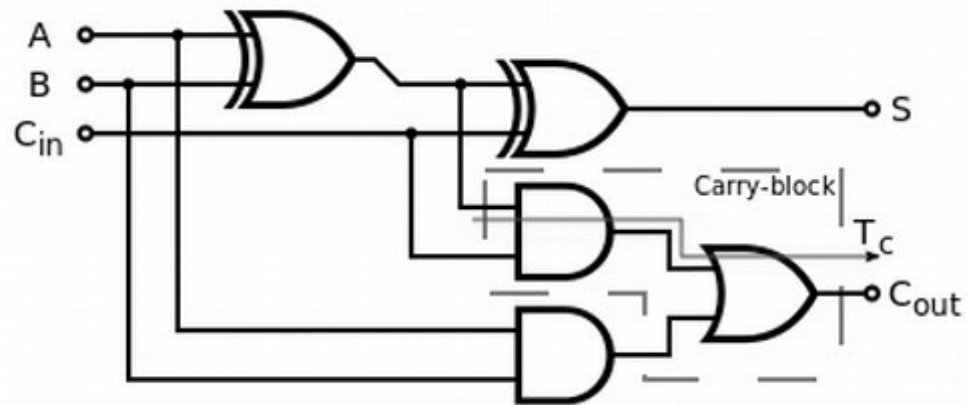
Adders and flip-flops



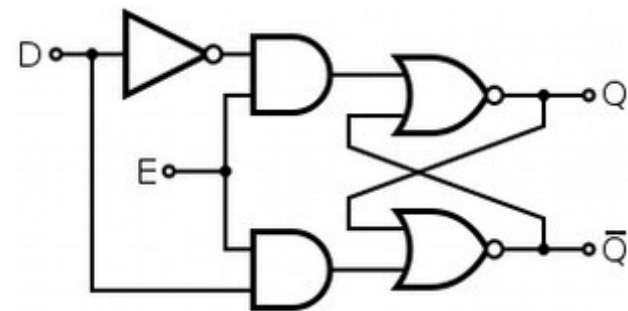
Half Adder



SR Flip-flop



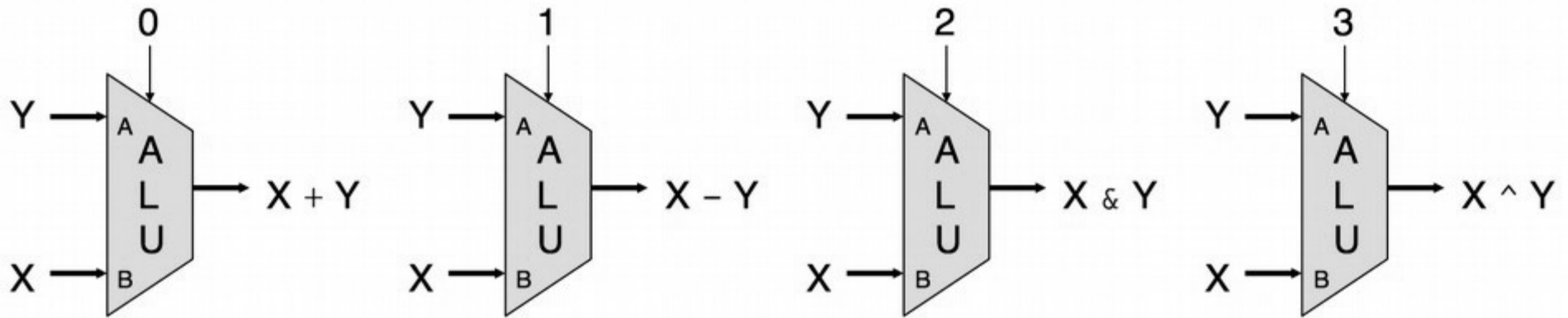
Full Adder



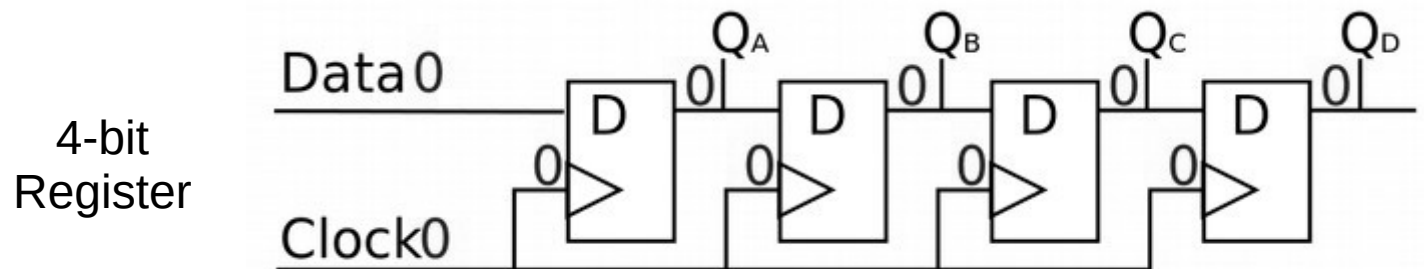
Gated D Flip-flop

ALUs and memory

- Combine **adders** and **multiplexors** to make **arithmetic/logic units**
- Combine **flip-flops** to make **register files** and **main memory**

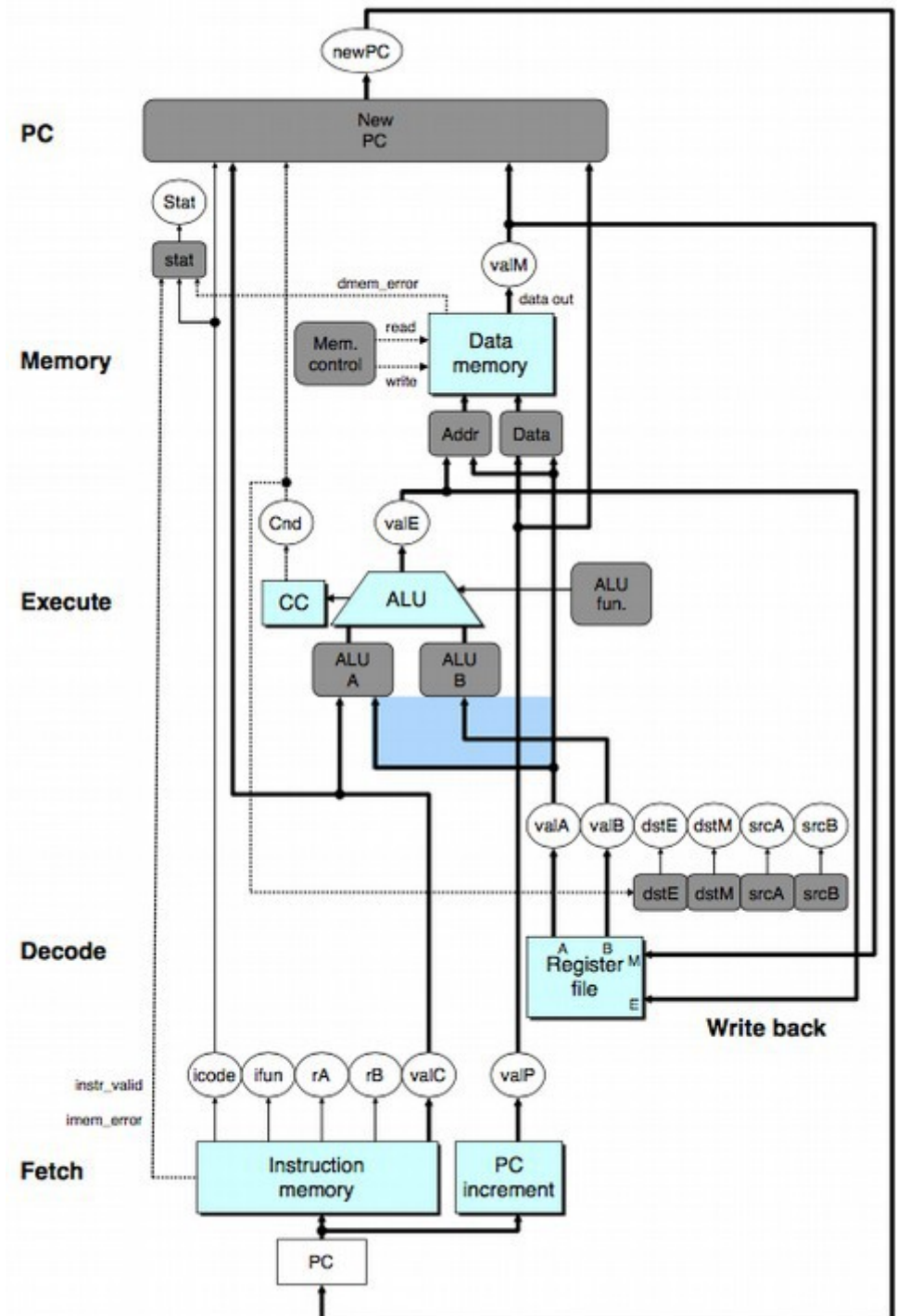


Basic Arithmetic Logic Unit (ALU)



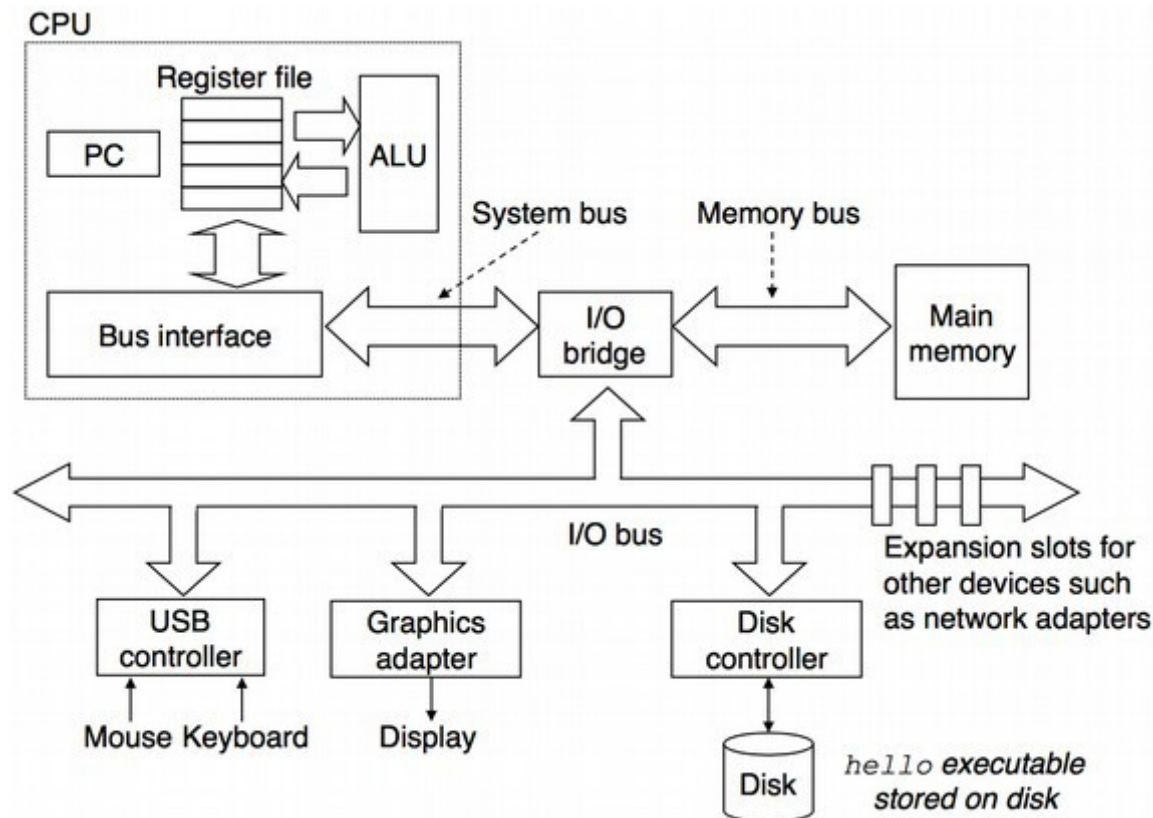
CPUs

- Combine **ALU** with **registers** and **memory** to make CPUs



Computers

- Combine **CPU** with other electronic components and devices (similarly constructed) communicating via **buses** to make a **computer**



Big picture

- Basic systems design approach: exploit **abstraction**
 - Start with simple components
 - Combine to make more complex components
 - Repeat using the new components as **black box** “simple components”
- This is true of most areas in systems
 - **CS 261**: transistors → gates → circuits → adders/flip-flops → ALUs/registers → CPUs/memory → computers
 - **CS 261**: machine code → assembly → C code → Java/Python code
 - **CS 361/470**: threads → processes → nodes → networks/clusters
 - **CS 432**: scanner → parser → analyzer → code generator → optimizer
 - **CS 450**: files + processes + I/O → kernel → operating system

Course status

- We've hit the bottom
 - Or at least as far down as we're going to go (logic gates)—from here we go back up!
- Next week
 - Combinational circuits
 - Sequential circuits
 - CPU architecture

Suggestion: download **Logisim** (already installed on lap machines) and play around with some circuits!