# CS 261 Fall 2016 

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## Binary Arithmetic

## Binary Arithmetic

- Topics
- Basic addition
- Overflow
- Multiplication


## Basic addition

- Binary and hex addition are fundamentally the same as decimal addition
- Add digit-by-digit, using a carry as necessary
- Result generally requires more bits than the two operands



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- Binary and hex addition are fundamentally the same as decimal addition
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Figure 2.21 Integer addition. With a 4-bit word size, the sum could require 5 bits.

- Unsigned addition is subject to overflow
- Caused by truncation to integer size

(assume a 16-bit integer)


Figure 2.23 Unsigned addition. With a 4-bit word size, addition is performed
modulo 16.

## Overflow

- Two's complement addition
- Works just like unsigned addition mechanically
- Subject to both positive and negative overflow
- Overflows if carry-in and carry-out differ for sign bit

| $x$ | $y$ | $x+y$ | $x+_{4}^{\mathrm{t}} y$ | Case |
| ---: | ---: | ---: | ---: | ---: |
| -8 | -5 | -13 | 3 | 1 |
| $[1000]$ | $[1011]$ | $[10011]$ | $[0011]$ |  |
| -8 | -8 | -16 | 0 | 1 |
| $[1000]$ | $[1000]$ | $[10000]$ | $[0000]$ |  |
| -8 | 5 |  | -3 | 2 |
| $[1000]$ | $[0101]$ |  | $[1101]$ |  |
| 2 | 5 |  | 7 | 3 |
| $[0010]$ | $[0101]$ |  | $[0111]$ |  |
| 5 | 5 | 10 | -6 | 4 |
| $[0101]$ | $[0101]$ | $[01010]$ | $[1010]$ |  |

Figure 2.24
Relation between integer and two's-complement addition. When $x+y$ is less than $-2^{w-1}$, there is a negative overflow. When it is greater than or equal to $2^{w-1}$, there is a positive overflow.


Figure 2.25 Two's-complement addition examples. The bit-level representation of the 4-bit two's-complement sum can be obtained by performing binary addition of the operands and truncating the result to 4 bits.

## Overflow

- Two's complement addition
- Works just like unsigned addition mechanically
- Subject to both positive and negative overflow
- Overflows if carry-in and carry-out differ for sign bit


Figure 2.26 Two's-complement addition. With a 4-bit word size, addition can have a negative overflow when $x+y<-8$ and a positive overflow when $x+y \geq 8$.

## Multiplication

- Like addition, fundamentally the same as base 10
- Actually, it's even simpler!
- Same regardless of encoding

| Mode | $x$ |  | y |  | $x \cdot y$ |  | Truncated $x \cdot y$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Unsigned | 5 | [101] | 3 | [011] | 15 | [001111] | 7 | [111] |
| Two's complement | -3 | [101] | 3 | [011] | -9 | [110111] | -1 | [111] |
| Unsigned | 4 | [100] | 7 | [111] | 28 | [011100] | 4 | [100] |
| Two's complement | -4 | [100] | -1 | [111] | 4 | [000100] | -4 | [100] |
| Unsigned | 3 | [011] | 3 | [011] | 9 | [001001] | 1 | [001] |
| Two's complement | 3 | [011] | 3 | [011] | 9 | [001001] | 1 | [001] |

Figure 2.27 Three-bit unsigned and two's-complement multiplication examples. Although the bit-level representations of the full products may differ, those of the truncated products are identical.

## Multiplication

- Special case: multiply by powers of 2 (shift left)

| $2 \ll 1=4$ | $(2 * 2)$ |  |
| :--- | :--- | :--- |
| $1 \ll 2=4$ | $(1 * 2 * 2)$ |  |
| $1<4=16$ | $(1 * 2 * 2 * 2 * 2)$ |  |
| $4 \ll 1=8$ | $(4 * 2)$ |  |
| $4<2=16$ | $(4 * 2 * 2)$ |  |

## - General case is expensive!

- Special case: divide by powers of two (shift right)

| k | $\gg \mathrm{k}$ (binary) | Decimal | $12,340 / 2^{\mathrm{k}}$ |
| :--- | :---: | :---: | :---: |
| 0 | 0011000000110100 | 12,340 | $12,340.0$ |
| 1 | 0001100000011010 | 6,170 | $6,170.0$ |
| 4 | 0000001100000011 | 771 | 771.25 |
| 8 | 0000000000110000 | 48 | 48.203125 |

Figure 2.28 Dividing unsigned numbers by powers of 2. The examples illustrate how performing a logical right shift by $k$ has the same effect as dividing by $2^{k}$ and then rounding toward zero.

| k | $\gg \mathrm{k}$ (binary) | Decimal | $-12,340 / 2^{\mathrm{k}}$ |  |
| :--- | :---: | ---: | :---: | :---: |
| 0 | 1100111111001100 | $-12,340$ | $-12,340.0$ |  |
| 1 | 1110011111100110 | $-6,170$ | $-6,170.0$ |  |
| 4 | 1111110011111100 | -772 | -771.25 | Two's complement |
| 8 | 1111111111001111 | -49 | -48.203125 |  |

Figure 2.29 Applying arithmetic right shift. The examples illustrate that arithmetic right shift is similar to division by a power of 2, except that it rounds down rather than toward zero.

## Division

- General case is expensive!
- Special case: divide by powers of two (shift right)

| k | Bias | $-12,340+$ bias (binary) | $\gg \mathrm{k}$ (binary) | Decimal | $-12,340 / 2^{\mathrm{k}}$ |
| :--- | ---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 1100111111001100 | 1100111111001100 | $-12,340$ | $-12,340.0$ |
| 1 | 1 | 1100111111001101 | 1110011111100110 | $-6,170$ | $-6,170.0$ |
| 4 | 15 | 1100111111011011 | 1111110011111101 | -771 | -771.25 |
| 8 | 255 | 1101000011001011 | 1111111111010000 | -48 | -48.203125 |

Figure 2.30 Dividing two's-complement numbers by powers of 2. By adding a bias before the right shift, the result is rounded toward zero.

## Quiz

- What is $5+2$ ?
-What is $4-3$ ?
-What is $2 \ll 3$ ?
-What is $3 \ll 3$ ?
-What is $16 \gg 2$ ?

Show your work using two's complement in both hex and binary using 8-bit integers

