

CS240

Fall 2015

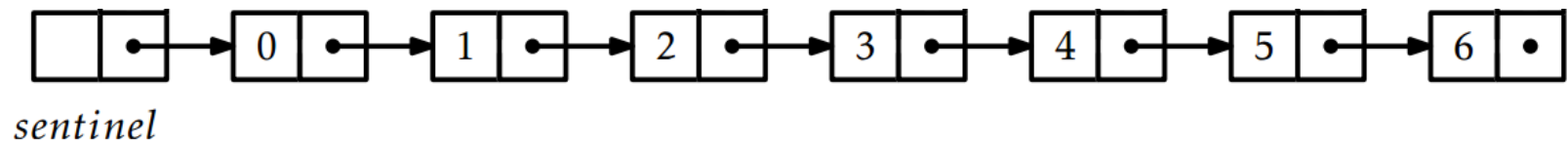
Mike Lam, Professor

Skip Lists

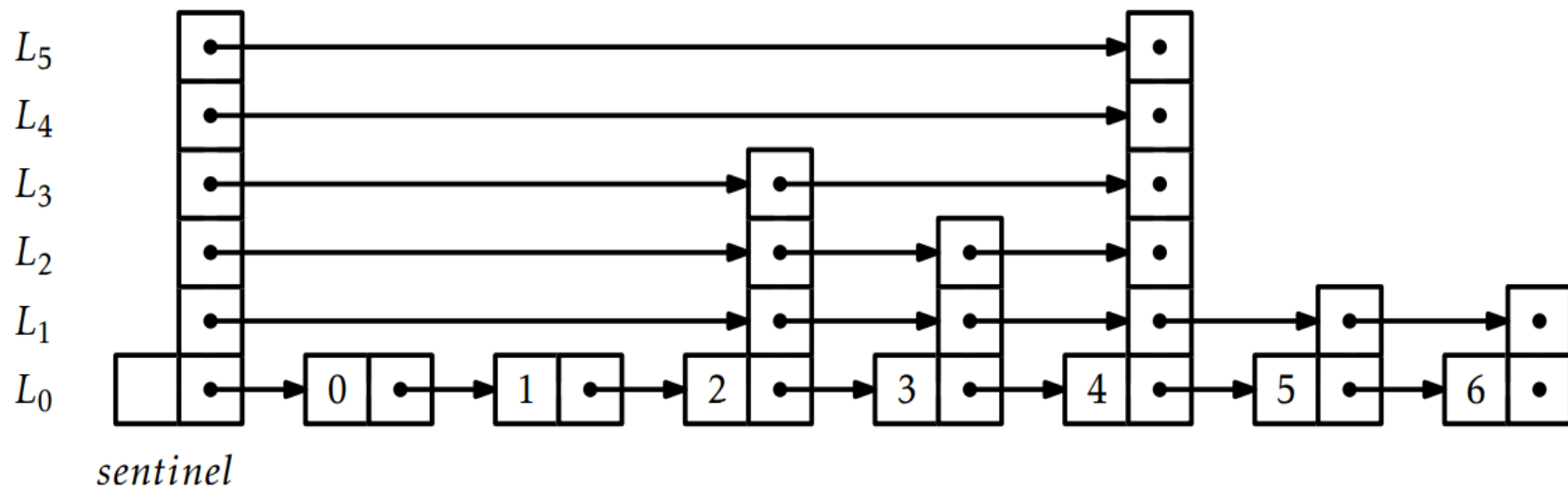
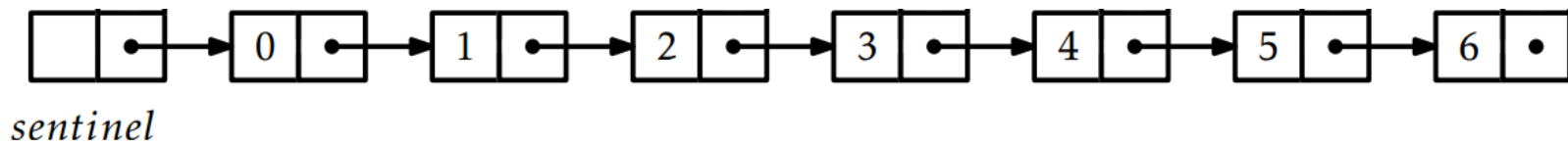
Skip Lists

- Multi-level sorted linked list
 - Multiple **next** pointers per node
- Higher levels contain “skip” links to locations further down the list
- All nodes reside in bottom level (L_0)
- Some nodes reside in higher levels as well
 - The height of the node is the number of levels in which that node resides
 - i.e., the height of a node is the number of **next** pointers
 - **Caution:** heights are 1-based while levels are 0-based

Skip Lists



Skip Lists



Skip List

find(x):

Basic idea: **move right** then **move down** (and repeat!)

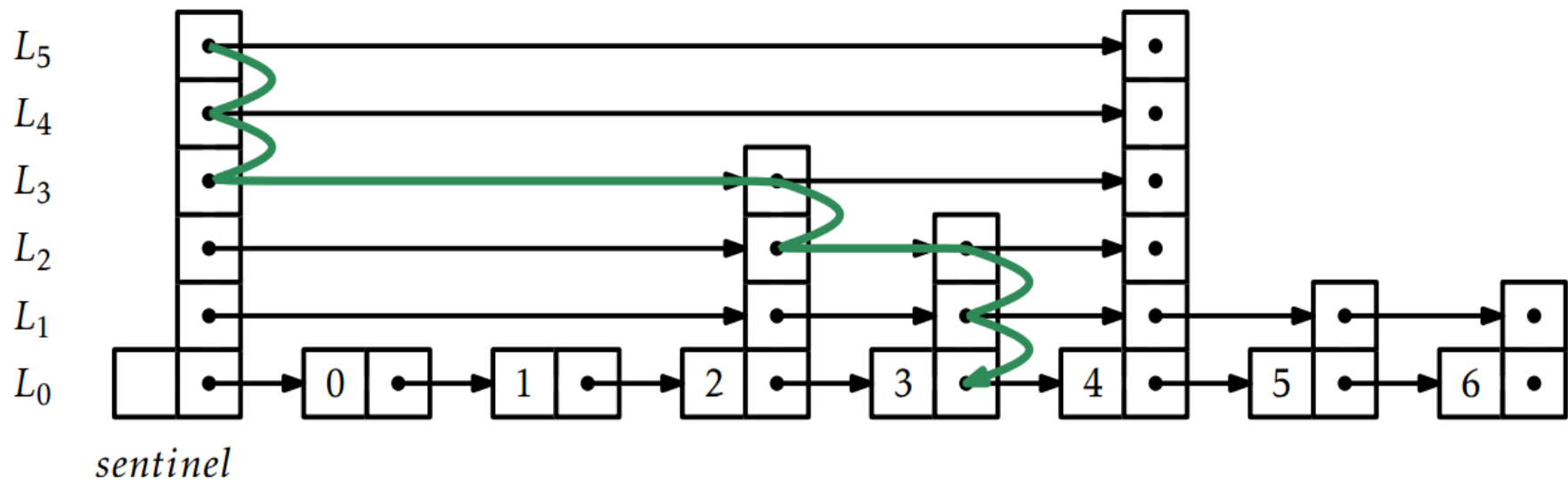
$u = \text{head}$, $r = \text{max_height} - 1$

while $r \geq 0$:

while $u.\text{next}[r] \neq \text{NULL}$ and $u.\text{next}[r].\text{data} \leq x$:

$u = u.\text{next}[r]$

$r = r - 1$

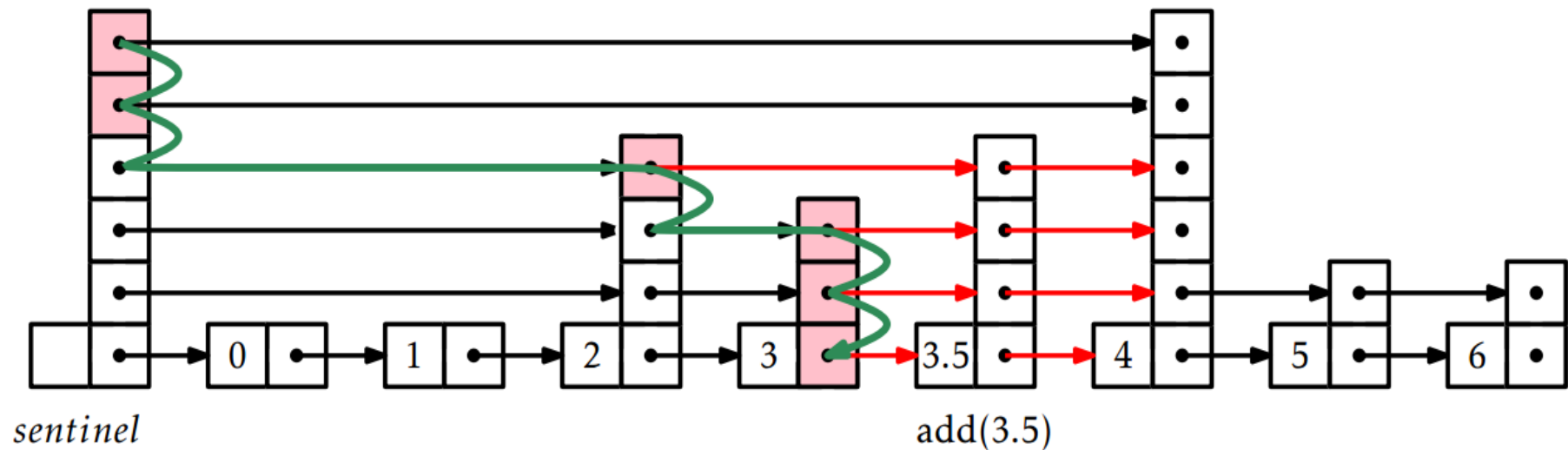


Adding to a Skip List

- Key concept: randomization
 - Flip a coin several times to determine the maximum height of a node
 - If we get heads, keep flipping!
 - Examples
 - Tails on first try: height = 1
 - Tails on second try: height = 2
 - Tails on third try: height = 3
 - etc.
 - Expected height of a node is 1
 - Consequence: half of all nodes only reside on bottom level

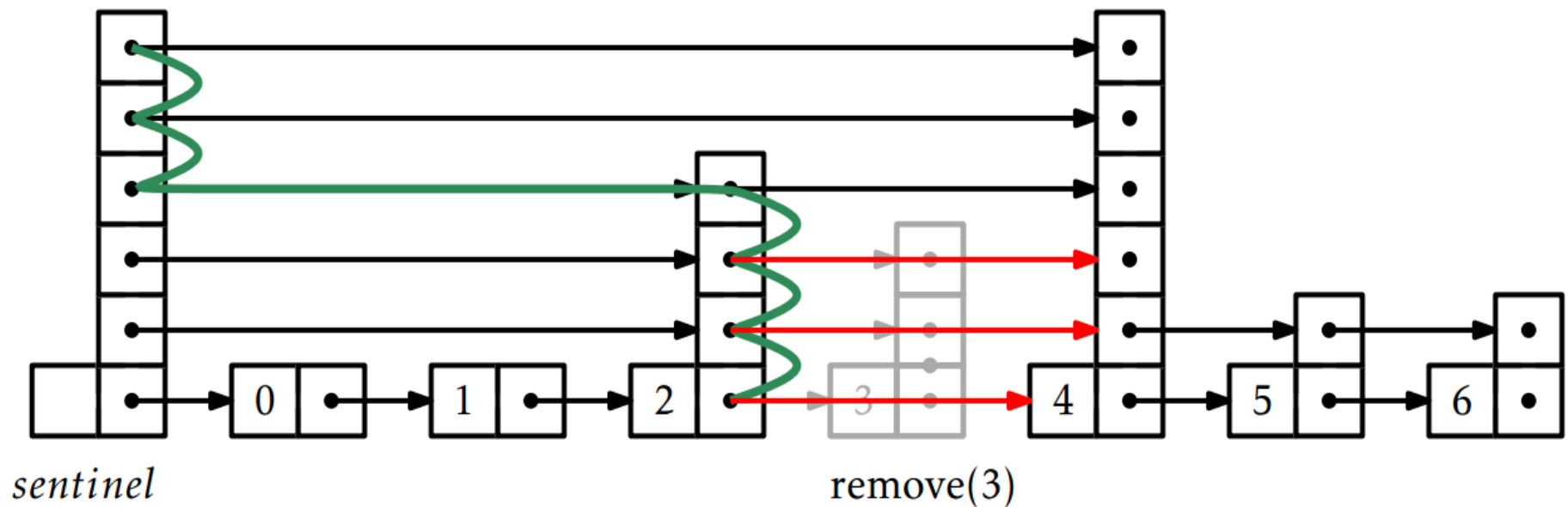
Adding to a Skip List

Basic idea: **move right** then **move down** (and repeat!)
Insert new node into all levels lower than its height



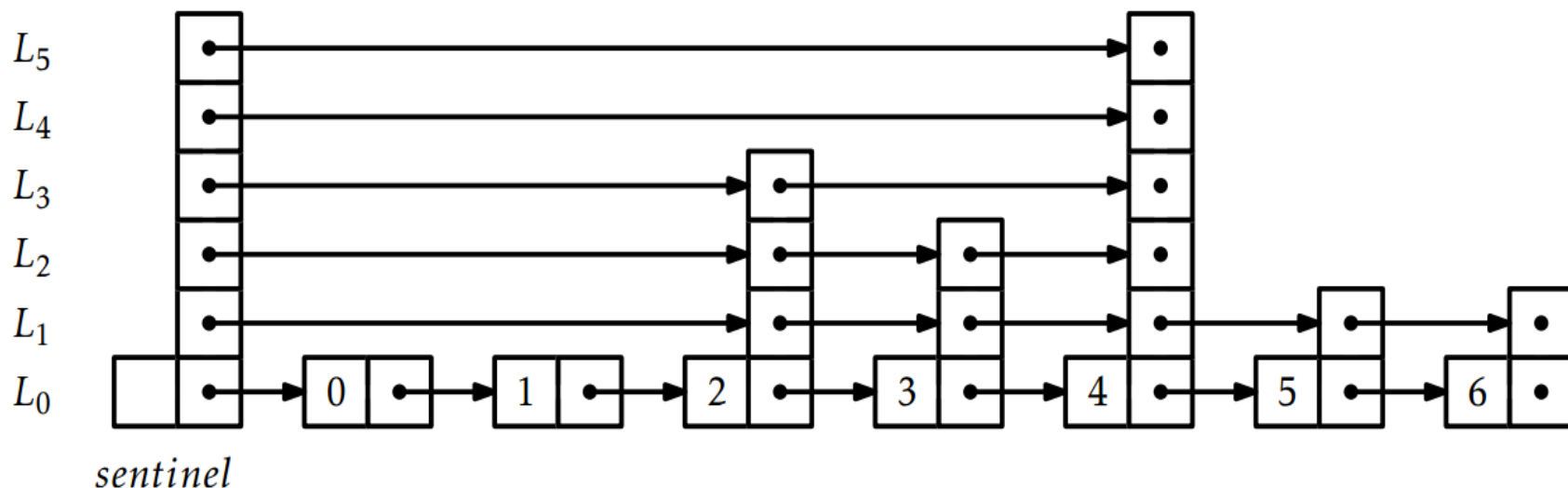
Removing from a Skip List

Basic idea: **move right** then **move down** (and repeat!)
Remove node from any level where it is present



Skip Lists

- Draw the results of the following operations:
 - **find**(1) and **find**(2) and **find**(5)
 - **add**(1.5) and **add**(4.5)
 - height = 3
 - **remove**(0) and **remove**(2)

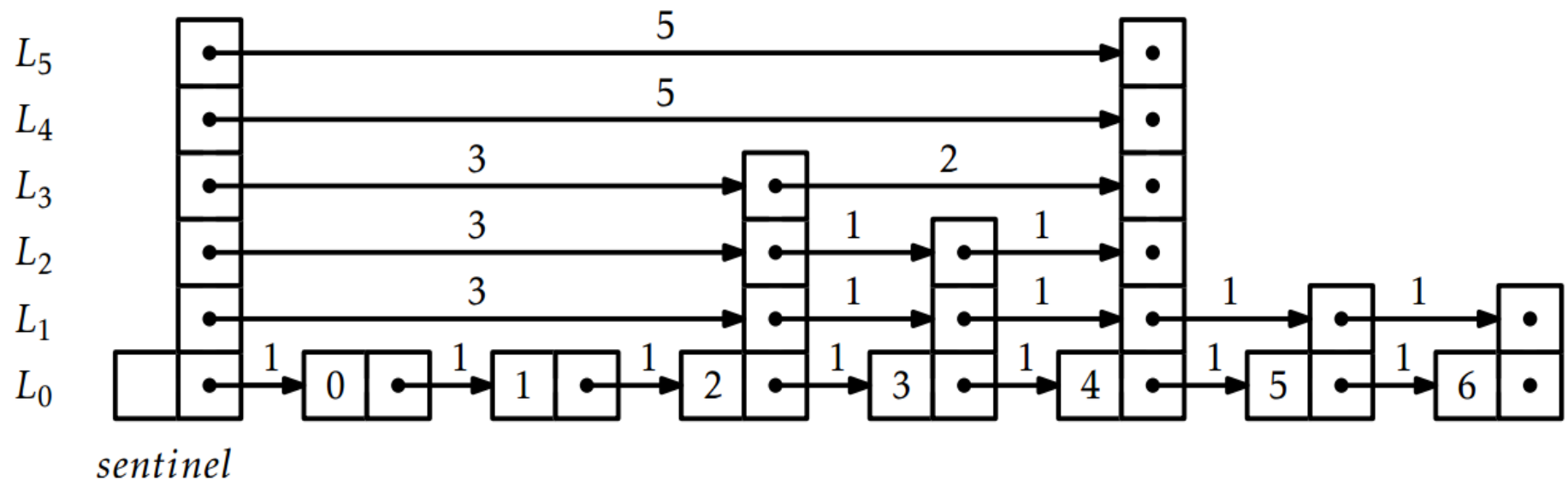


Skip Lists

- Insert, remove, and search all involve a traversal of the list
 - Move right as far as you can, then go down a level and repeat
 - Skips as many nodes as possible
- Visit an expected constant number of nodes per level
 - The constant is exactly 2 in a "perfect" skip list
- There will be approximately $\log n$ levels
 - 1/2 of nodes are only on bottom level, 1/4 are in bottom two levels, 1/8 are in bottom three levels, etc.
 - Halving of input space at each level
- Thus, all three operations are $O(\log n)$!

Link Lengths

- Annotate a skip list with *length* info on links
 - Each length is the sum of the lengths of links below it



Now we can find the item at a given index in $O(\log n)$ as well!

Updating Link Lengths

