

CS240 Fall 2014

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recursion

n. [ri-**kur**-zhuh n]

1. See “recursion”

Recursion

Recursion

- The expression of a problem solution in a way that depends on solutions to smaller instances of the same problem
- For some problems, a recursive solution is cleaner than the corresponding iterative solution
- Classics:
 - A **list** is either 1) an “empty list” or 2) an item followed by a **list**
 - **fact**(n) = 1 if $n \leq 1$, $n * \mathbf{fact}(n-1)$ if $n > 1$
 - Tower of Hanoi / Brahma

Recursion

- The language runtime handles the actual semantics of recursive behavior
- Usually, it tracks recursive calls using a stack
- Every function call pushes a new entry (called an “activation record” or “frame”) to the stack
- A record is popped when a function returns, and execution returns to the function on the top of the stack

Recursion

- “Call stack”
- Details are machine- and language-dependent
- More info in CS430

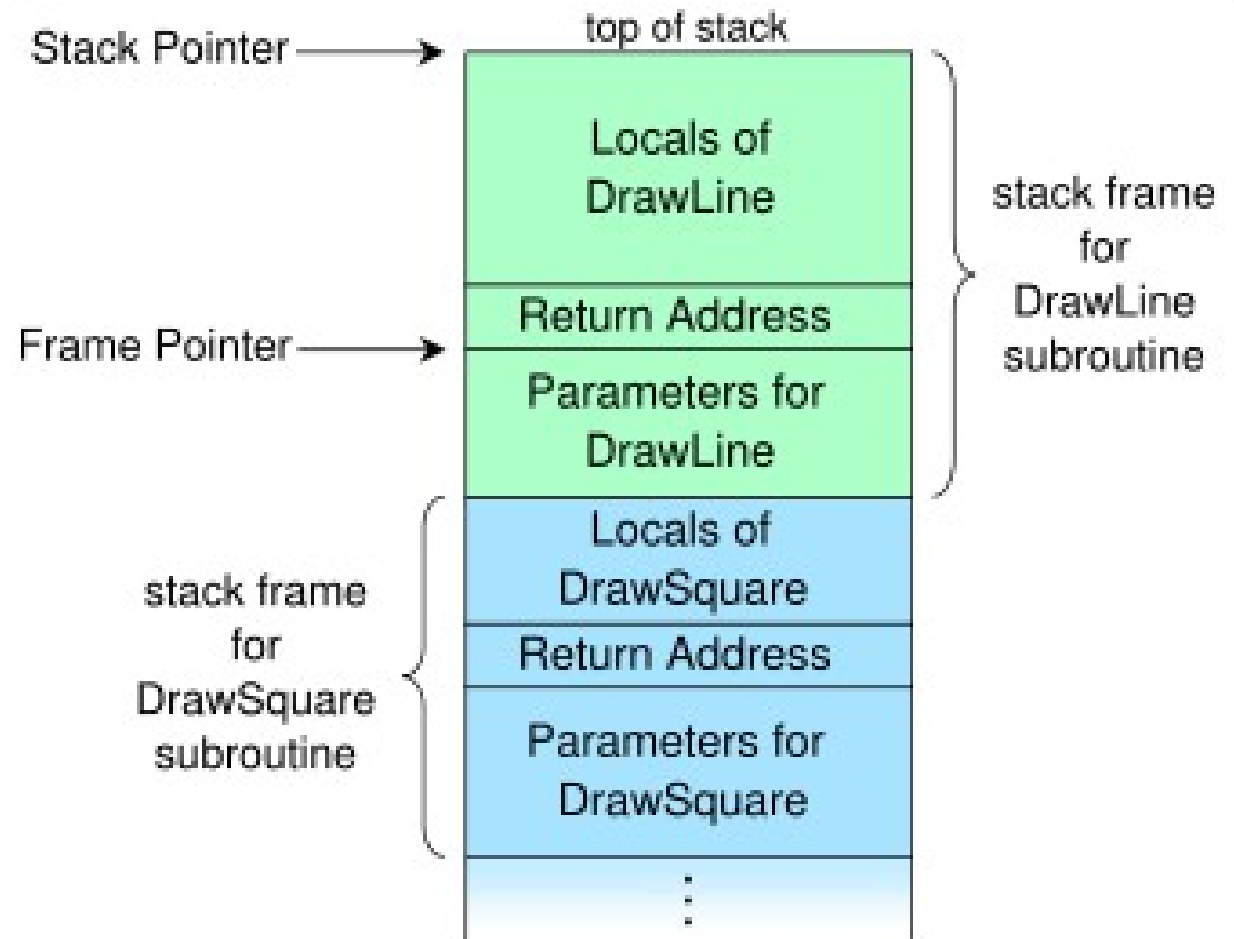


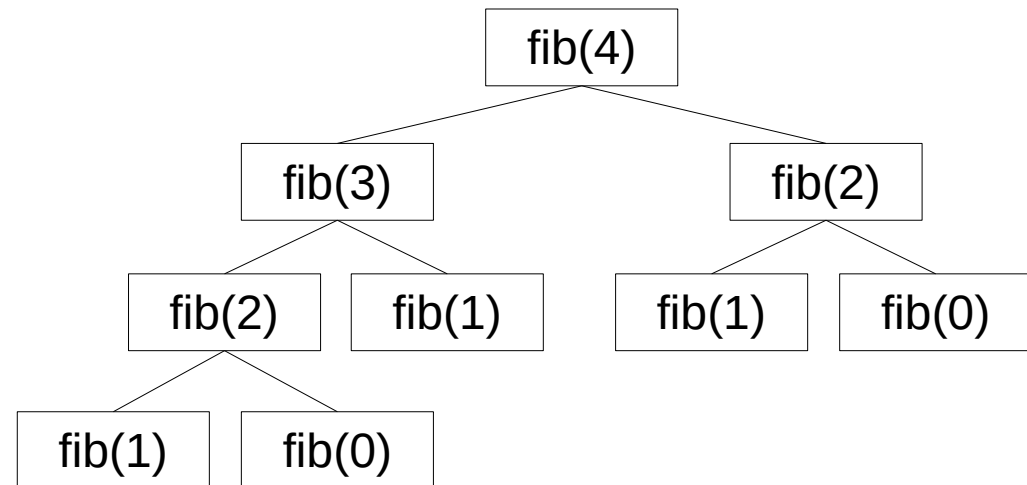
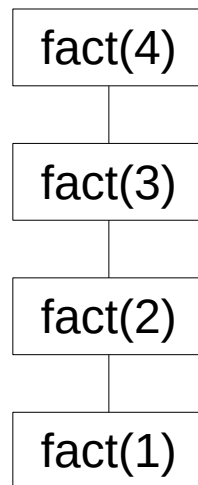
Image from Wikipedia article “Call Stack”

Recursion

- Single vs. binary vs. multiple recursion
 - **fact**(n) = 1 if $n \leq 1$, $n * \mathbf{fact}(n-1)$ if $n > 1$
 - **fib**(n) = 1 if $n \leq 1$, $\mathbf{fib}(n-1) + \mathbf{fib}(n-2)$ if $n > 1$
- Trace: fact(4) vs. fib(4)

Recursion

- Single vs. binary vs. multiple recursion
 - **fact**(n) = 1 if $n \leq 1$, $n * \mathbf{fact}(n-1)$ if $n > 1$
 - **fib**(n) = 1 if $n \leq 1$, $\mathbf{fib}(n-1) + \mathbf{fib}(n-2)$ if $n > 1$
- Trace: **fact**(4) vs. **fib**(4)



Binary Search

```
def find(array, item):
    return helper(array, item, 0, len(array))

def helper(array, item, left, right):
    mid = (right-left)//2 + left
    if array[mid] > item:
        return helper(array, item, left, mid)
    elif array[mid] < item:
        return helper(array, item, mid+1, right)
    else:
        return mid < len(array) and array[mid] == item
```

Binary Search

- Trace: `find([1,4,5,7,9,11,15], 5)`

Binary Search

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left = 0
right = 7

Binary Search

- Trace: `find([1,4,5,7,9,11,15], 5)`

left = 0
right = 7

mid = 3

[1, 4, 5, 7, 9, 11, 15]

Binary Search

- Trace: `find([1,4,5,7,9,11,15], 5)`

left = 0
right = 7

mid = 3

[1, 4, 5, 7, 9, 11, 15]

[1, 4, 5, 7, 9, 11, 15]

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Binary Search

- Trace: `find([1,4,5,7,9,11,15], 5)`

left = 0
right = 7

mid = 3 [1, 4, 5, 7, 9, 11, 15]

[1, 4, 5, 7, 9, 11, 15]

left = 0
right = 3

mid = 1 [1, 4, 5, 7, 9, 11, 15]

Binary Search

- Trace: `find([1,4,5,7,9,11,15], 5)`

left = 0
right = 7

mid = 3 [1, 4, 5, 7, 9, 11, 15]

[1, 4, 5, 7, 9, 11, 15]

left = 0
right = 3

mid = 1 [1, 4, 5, 7, 9, 11, 15]

[1, 4, 5, 7, 9, 11, 15]

left = 2
right = 3

Binary Search

- Trace: `find([1,4,5,7,9,11,15], 5)`

left = 0
right = 7

mid = 3 [1, 4, 5, 7, 9, 11, 15]

[1, 4, 5, 7, 9, 11, 15]

left = 0
right = 3

mid = 1 [1, 4, 5, 7, 9, 11, 15]

[1, 4, 5, 7, 9, 11, 15]

left = 2
right = 3

mid = 2 [1, 4, 5, 7, 9, 11, 15]

A note on the word “binary”

- Binary search is not binary recursion!
 - Only recurses on one half of the list
 - So it is single recursion
- Binary sum is binary recursion
 - Recurses on both sides of the list

Binary Search

- What is the running time of a binary search?

Binary Search

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- Need a way to express recursion costs mathematically
- Write a function!
 - Express $T(n)$ in terms of itself

Binary Search

- What is the running time of a binary search?
- Need a way to express recursion mathematically
- Write a function!
 - Express $T(n)$ in terms of itself
- For binary search: $T(n) = 1 + T(n/2)$
 - To search n items, do one comparison then recurse on the appropriate half-list

Recurrences

- Recursive formulas are called “recurrences”
- We still want to find a “closed-form” description
 - Something like 2^n or “log n” or 5^n
- We will talk more on Wednesday about how to solve recurrences
- But first, we need to be comfortable tracing recursive code

Exercise 1

- Given the following code:

```
def foo(n):  
    if n < 2:  
        return 1  
    else:  
        return n * foo(n-1)
```

- Trace the following call:

```
print(str(foo(4)))
```

Exercise 2

- Given the following code:

```
def bar(text):  
    if len(text) <= 1:  
        return True  
    return text[0] == text[-1] and  
           bar(text[1:-1])
```

- Trace the following call:

```
print(str(bar("abbaba")))
```

Exercise 3

- Given the following code:

```
def baz(x, n):  
    if n == 0:  
        return 1  
    y = baz(x, n//2)  
    if n % 2 == 1:  
        return x * y * y  
    else:  
        return y * y
```

- Trace the following call:

```
print(str(baz(2, 10)))
```

Exercise 4

- Given the following code:

```
def hanoi(n, src, dst, tmp):  
    if n == 1:  
        print("move from " + str(src) +  
              " to " + str(dst))  
    else:  
        hanoi(n-1, src, tmp, dst)  
        hanoi( 1, src, dst, tmp)  
        hanoi(n-1, tmp, dst, src)
```

- Trace the following call:

```
hanoi(3, "a", "c", "b")
```